# artificial intelligence on the dragon computer 

 make your micro think
## keith and steven brain



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## CONTENTS

Page
Introduction ..... V11
1 Artificial Intelligence ..... 1
2 Just Following Orders ..... 5
3 Understanding Natural Language ..... 21
4 Making Reply ..... 37
5 Expert Systems ..... 55
6 Making Your Expert System Learn for Itself 71
7 Fuzzy Matching ..... 85
8 Recognising Shapes ..... 97
9 An Intelligent Teacher ..... 109
10 Putting It All Together ..... 115

## Contents in detail

## CHAPTER I

Artificial Intelligence
Fantasy - reality: two-way conversations. robots, expert systems.

## CHAPTER 2

Just Following Orders
Preset orders and fixed responses - DATA arrays - expanding the vocabulary - removing redundancy - abbreviated commands - partial matching - sequential commands.

## CHAPTER 3

Understanding Natural Language
Dealing with sentences - subjects, objects. verbs, adjectives. adverbs punctuation - a sliding search - rearranging the word store array.

## CHAPTER 4

Making Reply
Getting more sensible replies - making logical decisions before replying choosing the correct subject - problems with objects - changing tense.

## CHAPTER 5

Expert Systems
How an expert works - simple problems - more difficult problems including pointers - sequential and parallel branching - checking how well the answers match the data - better in bits.

CHAPTER 6
Making Your Expert System Learn for Itself
Letting the computer work out its own rules for two objects - a wider spectrum - watching what happens.

## CHAPTER 7

Fuzzy Matching
Recovering information from the human mind - Soundex Coding - a computer program for converting names - retrieving information.

## CHAPTER 8

Recognising Shapes
Simulating the action of a light sensor - inserting into sentences - a branching short cut.

## CHAPTER 9

An Intelligent Teacher
Questions and answers - keeping a score - shifting the emphasis of questions to areas of difficulty - making questions easier or harder.

## CHAPTER 10

Putting It All Together
Making conversation with the computer - making decisions, cost arrays and profit arrays - the Computer Salesman.

## Introduction

Artificial Intelligence is undoubtedly an increasingly important area in computer development which will have profound effects on all our lives in the next few decades. The main aim of this book is to introduce the reader to some of the concepts involved in Artificial Intelligence and to show them how to develop 'intelligent' routines in BASIC which they can then incorporate into their own particular programs. Only a superficial knowledge of BASIC is assumed, and the book works from first principles as we believe that this is essential if you are really to understand the problems involved in producing intelligence, and how to set about overcoming them.

The basic format of the book is that ideas are taken and suitable routines built up step by step, exploring and comparing alternative possibilities wherever feasible. Rather than simply giving you a series of completed programs, we encourage you to experiment with different approaches to let you see the results for yourself. Detailed flowcharts of most of the routines are included. The main emphasis in the routines is placed on the Al aspects and we have therefore avoided 'tarting up' the screen display as this tends to obscure the significance of the program. In places you may notice that odd lines are redundant. but these have been deliberately included in the interests of clarity of program flow. As far as possible, retyping of lines is strenuously avoided but modification of lines is commonplace. All listings in the book are formatted so that they appear as you will see them on the screen. In most cases, spaces and brackets have been used liberally to make listings easier to read but be warned that some spaces and brackets are essential so do not be tempted to remove them all. All routines have been rigorously tested and the listings have been checked very thoroughly so we hope that you will not find any bugs. It is a sad fact of lifethat most bugs arise as a result of 'tryping mitsakes' by the user. Semi-colons and commas may look very insignificant but their absence can have very profound effects!

Artificial Intelligence is increasing in importance every day and we hope that this book will give you a usef ul insight into the area. Who knows - if you really work at the subject you might be able to persuade your machine to read our next book for itself!

## CHAPTER 1

## Artificial Intelligence

## Fantasy

For generations, science fiction writers have envisaged the development of intelligent machines which could carry out many of the functions of man himself (or even surpass him in some areas), and the public image of Artificial Intelligence has undoubtedly been coloured by these images. The most common view of a robot is that it is an intelligent machine of generally anthropomorphic (human) form which is capable of independently carrying out instructions which are given to it in only a very general manner.

Of course, most people have ingrained Luddite tendencies when it comes to technology so in the early stories these robots tended to have a very bad press, being cast in the traditional role of the 'bad guys' but with nearinvincibility and lack of conscience built in. The far-sighted Isaac Asimov wove a lengthy series of stories around his concept of 'positronic robots' and was probably the first author really to get to grips with the realities of the situation. He laid down his famous 'Three Laws of Robotics' which specified the basic ground rules which must be built into any machine which is capable of independent action - but it is interesting to note that he could not foresee the time when the human race would accept the presence of such robots on the earth itself.
'Star Wars' introduced the specialised robots R2D2 and C3PO, but we feel that many of their design features were a little strange. Perhaps there is an Interplanetary Union of Robots, and a demarcation dispute prevented direct communication between humans and R2D2. In 'The Stepford Wives', the local husbands got together and had the (good?) idea of converting their wives into androids who automatically did exactly what was expected of them, but the sequel revealed the dangers of the necessity to continuously reinforce with an external stimulus! Perhaps one hope for mankind is that any aliens who chance upon us will not have watched 'Battlestar Galactica', and will therefore build robots of the Cylon type who, rather like the old Space Invaders, are always eventually defeated because they are totally predictable.

Of course intelligent computers also appear in boxes without arms and legs, although flashing lights seem obligatory. Input/output must obviously be vocal but the old metallic voice has clearly gone out of fashion
in favour of some more definite personality. If all the boxes look the same then this must be a good idea, but please don't make yours all sound like Sergeant-Major Zero from 'Terrahawks'! Michael Knight's KITT sounds like a reasonable sort of machine to converse with, and it is certainly preferable to the oily SLAVE and obnoxious ORAC from 'Blake's Seven'. ORAC seemed to pack an enormous amount of scorn into that little perspex box, but other writers have appreciated the difficulties which may be produced if you make the personality of the machine too close to that of man himself.

In Arthur C. Clarke's '2001: A Space Odyssey', the ultimately-intelligent computer HAL eventually had a nervous breakdown when he faced too many responsibilities; but in ‘Dark Star’ the intelligent bomb was quite happy to discuss Existentialism with Captain Doolittle but was unwilling to deviate from his planned detonation time, although still stuck in the bomb bay. In 'The Restaurant At The End of The Universe', the value of the Sirius Cybernetics Corporation Happy Vertical People Transporter was reduced significantly when it refused to go up as it could see into the future and realised that if it did so it was likely to get zapped; and the NutriMatic Drinks Synthesiser was obviously designed by British Rail Catering as it always produced a drink that was 'almost, but not quite, entirely unlike tea'.

More worrying themes have also recently appeared. The most significant feature of 'Wargames' was not that someone tapped into JOSH UA (the US Defence Computer), but that once the machine started playing thermonuclear war it wouldn't stop until someone had won the game. And in 'The Forbin Project' the US and Russian computers got together and decided that humans are pretty irrelevant anyway. Of course, if you are Marvin the Paranoid Android and have a brain the size of a planet and a Genuine People Personality, you can succeed without weapons by confusing the enemy machine into shooting the floor from under itself whilst discussing your personal problems.

## Reality

The definition and recognition of machine intelligence is the subject of fast and furious debate amongst the experts in the subject. The most generallyaccepted definition is that first proposed by Alan Turing way back in the late 1940 s when computers were the size of houses and even rarer than a slide-rule is today. Rather than trying to lay down a series of criteria which must be satisfied, he took a much broader view of the problem. He reasoned that most human beings accept that most other human beingsare intelligent and that therefore if a man cannot determine whether he is dealing with another man (or woman), or only with a computer, then he must accept that such a machine is intelligent. This forms the basis of the
famous 'Turing Test', in which an operator has to hold a two-way conversation with another entity via a keyboard and try to get the other party to reveal whether it is actually a machine or just another human being - very awkward!

Many fictional stories circulate about this test, but our favourite is the one where a job applicant is set down in front of a keyboard and left to carry on by himself. Of course he realises the importance of this test to his career prospects and so he struggles valiantly to find the secret, apparently without success. However after some time the interviewer returns, shakes him by the hand, and congratulates him with the words 'Well done, old man, the machine couldn't tell if you were human so you are just what we need as one of Her Majesty's Tax Inspectors!'

Everyone has seen from TV advertisements that the use of computeraided design techniques is now very common, and that industrial robots are almost the sole inhabitants of car production lines (leading to the car window sticker which claims 'Designed by a computer, built by a robot, and driven by an idiot'). In fact, most of these industrial robots are really of minimal intelligence as they simply follow a pre-defined pathway without making very much in the way of actual decisions. Even the impressive paint-spraying robot which faithfully follows the pattern it learns when a human operator manually moves its arm cannot learn to deal with a new object without further human intervention.

On the other hand, the coming generation of robots have moresophisticated sensors and software, which allow them to determine the shape, colour, and texture of objects, and to make more rational decisions. Anyone who has seen reports of the legendary 'Micromouse' contests, where definitely non-furry electric vermin scurry independently and purposefully (?) to the centre of a maze, will not be a MAZEd by our faith in the future of the intelligent robot, although there seems little point in giving it two arms and two legs.

Another important area where Artificial Intelligence is currently being exploited is in the field of expert systems, many of which can do as well (or even better) than human experts, especially if you are thinking about weather forecasting. These systems can be experts on any number of things but, in particular, they are of increasingimportance in medical diagnosis and treatment - although the medical profession doesn't have to worry too much as there will always be a place for them since 'computers can't cuddle'

A major barrier to the wider use of computers is the ignorance and pigheadedness of the users, who will only read the instructions as a last resort, and who expect the machine to be able to understand all their little pecularities. Processing of 'natural language' is therefore a major growth area and the 'fifth generation' of computers will be much more userfriendly.

Most of the serious work on Artificial Intelligence uses more suitable (but exotic) languages than BASIC, such as LISP and PROLOG, which are pretty unintelligible to the average user and are probably not available for your home micro in any case. The BASICroutines which follow cannot therefore be expected to give you the key to world domination, although they should give you a reasonable appreciation of the possibilities and problems which Artificial Intelligence brings.

## CHAPTER 2 <br> Just Following Orders

As your computer is actually totally unintelligent，you can only converse with it in very simple terms．The first step，used in many simple adventure games，is to have a series of preset orders to which there are fixed responses． Let＇s start by taking a look at giving compass directions for which way to move．At first sight，the simplest way to program this appears to be to ask for an INPUT from the user and to write a separate IF－THEN line for each possibility（see Flowchart 2．1）．

1F3 FRINT＂DIPECTIOFた＂；
130 INPUT IN\＄
 DPTH＂
31日 TF 【性 $=" \mathrm{SOUTH}$ THER FPIr•1T＂ OUTH＂
20日 IF IH客＂＂以EST＂THEN FRIYT＂WE ST＂
230［F If 㨁＂EFHCT＂THEH PRIFTT＂EA ST＂
：35ด GOTD 109


Flowchart 2．1 Giving Compass Directions

If you type in anything other than the four key command words, nothing will be printed except for another input request. It would be more userfriendly if the computer indicated more clearly that this command was not valid. You could do that by including a test which shows that none of the command words has been found, but this becomes very long-winded, and effectively impossible when you have a long list of valid words.

 FSST" THEN PPINT"INVFLIS REDUEST"

On the other hand, adding GOTO 100 to the end of each IF-THEN line will force a direct jump back to the INPUT when a valid command is detected. If all the IF tests are not true then the program falls through to line 240 which prints a warning. Making direct jumps back when a valid word is found is a good idea anyway, as it saves the system making unnecessary tests when the answer has already been found (see Flowchart 2.2).


Flowchart 2.2 Deleting Unnecessary Tests

```
2MM IF IH绍="HORTH" THEN FRINT- "N
OFTH" GOTO 10⿴
```

```
?1@ IF IN&="SOUTH" THEN PRINT "S
0UTH" : GOTO 1GET
200 IF IN&="WEST" THEN FRINT "WE
ST":GOTO 10口
F3G IF IN%:="ERST" THEN PRINT "EF
```



```
24g FPINT"INURLIO REDIIEST"
```

That will echo the command given on the screen but of courseit does not actually DO anything．As a model to work with，we willstart at a position defined as $\mathrm{X}=$ and $\mathrm{Y}=0$ and indicate movements as plus and minus in relation to this point．Notice that integer variables are used wherever possible，as they are processed faster than real numbers，and this also removes the possibility of clashing with reserved variables．

## $10 \mathrm{X}=\mathrm{m}$

We now need to add the real response to the command，as well as the message indicating that it has been understood（see Flowchart 2．3）．

```
2@ด IF IN&-"NORTH" THEN FRINT "H
ORTH" '准-1:COTG 100
21] IF IN海="SDUTH" THEN PRINT "'S
กIITH':'Y'=Y+1: \, %TTO 100
2? IF IN&="WEST" THEN PRINT "WE
ST":Yx%-1:GOTO 10N
30 IF IN&="EAST" THEN PRINT "EA
```



That modification actually shows your position appropriately，relative to the origin．So that you can see what is happening，and where you are，add a printout of your current position：

```
11ू PRINT"אי":兑,"Y":Y
```


## Using subroutines

Of course，that was a very simple example and，particularly where the results of your actions are more complicated，it is usually better to put the responses into subroutines．


219 IF IN＝＂SDIJTH＂THEN GOSIE 21明：GOTO 10日
2．2ด IF INक＝＂WEST＂THEN GDSUB 220日：GOTO 100
2：30 IF IN $=$＝＂ERST＂THEN GDSUB 230日：GOTO 1のロ

2月g PRINT＂GOING NDRTH＂：$\because=Y-1: R$ ETIUFN
？1旬 PRINT＂GOINE SDIUTH＂：$Y:=1 Y+1: R$ ETIJRN
22GGE PRINT＂GOING WEST＂：X＝X－1：RE TIJRN
230 PRINT＂GOINE ERST＂$X=X+1$ ：RE TI．JRN


Flowchart 2．3 Adding a Response

## More versatility

You could extend this use of IF THEN tests ad infinitum（or rather ad memoriam finitum！），but it is really a rather crude way of doing things which creates problems when you want to make your programs more sophisticated．A more versatile way to deal with command words and responses is to enter them as DATA and then store them in string arrays． First you must DI Mension arrays of suitable length for command words （C\＄）and responses（R\＄）．As variable－length strings are allowed（up to 255 characters）the actual text can be of almost any length．

## 30 DIM Cक（3），R甲（3）

If you put the commands and responses in pairs in the DATA statement， then it is more difficult to get them jumbled up and easier to read them in turn into the equivalent element in each array（see Table 2．1）．
 TH，GIINE SCUITH，LEST，TOING HEST，E RST，GORHE ERST
11 B0の FCJP H＝ด TO 3
11015 PERC Cक
1102 ANET H

| ELEMENT <br> NUMBER | COMMAND <br> WORD C\＄（n） | RESPONSE <br> R\＄（n） |
| :---: | :---: | :---: |
| 1 | NORTH | GOING NORTH |
| 2 | SOUTH | GOING SOUTH |
| 3 | WEST | GOING WEST |
| 4 | EAST | GOING EAST |

Table 2．1 Content of Command and Response Arrays

To initialise the arrays（fill them with your words），when you RUN add a GOSUB and RETURN．

All those IF-THEN tests can now be replaced by a single loop which compares your INPUT with each element of the array containing the command words (C\$) in turn (see Flowchart 2.4). Lines 200-220 need to be replaced by the following lines but notice also that line 230 must be deleted.

2GG FOR N二G TO 3

>: GOTD 1 คD
22G NEKT N


Flowchart 2.4 More Versatility

Now, IF your input, IN\$, corresponds to any of the command words, the program jumps out of the loop after printing the appropriate response, $\mathrm{R} \$(\mathrm{~N})$.

Of course we are now back in our original position of actually doing nothing, so we need to be able to call those action subroutines. First of all
let＇s arrange to jump out of the loop，if a match is found，t o a new routine at line 300.

## 210 IF $\mathrm{IN} \ddagger=[\$(\mathrm{~N})$ THEN PRINT P\＄（N ）：

We still have a pointer to indicate which word matched the input，as $N(t h e$ number of array elements checked）holds this value．We can use this in an ON－GOSUB lineto move to appropriate routines which are similar to the ones we wrote earlier，except that there is no need to def ine the particular message：this has already been printed as $\mathrm{R} \$(\mathrm{~N})$ ．

```
30D OH (N+1) GOSUE 2900.2100.220
```


こดดด ' $\%=Y-1$ : PETURN
? 1月の $Y=Y+1:$ RETURN
22日g $\mathrm{x}=\mathrm{K}-1$ : RETIURN
?ดด : $\because+1$ RETURN

## Expanding the vocabulary

The arrays can easily be expanded to contain more words．It would be better if we defined the number of words as a variable WD，which we would then use to DIMension the arrays and for both the filling and scanning loops．This produces a general routine which is easily modified．

```
20 WD=%
30 DIM Co( WD), RE(WIN)
30G FOR N*OTO W[P
1100D FOR N=8 TD WD
```

For example we can add intermediate compass directions which change both X and Y axes．

> 2ロ WDien？ 10010 DRTA NORTH ERST，GOING NORT H EAST，SOUTH ERST，TOINT，SOUTH EA ST，SOITH WEST，GOING SOUYH WEST，N OIPTH WEST，IOING NORTH WEST

and add some more subroutines：

300 ON（N＋1）GOSUB 290日，ช109．220 0，230日，2400，2500，260日，279ด：GOTח 1 日月

```
2400 Y Y=Y-1:X=\X41 : RETURN
2500 Y=Y+1:Y=X+1:RETI.JRN
2EgD Y=Y+1:X=X-1 :RETIJRN
270] Ya`\-1:K=X-1:RETIJRN
```


## Removing redundancy

All the responses so far haveincluded the word 'GOING' and th is word has actually been typed into each DATA statement. Now typing practice is very good for the soul but it would be much more sensible to define this common word as a string variable. Notice that a space is included at the end to space it from the following word.

## 

You can then delete all occurrences of this word in the DATA and combine G\$ with each key word in the response instead.

```
210 IF IN%=C@(M) THEN PRINT G%;R
S(N ):GOTO 300
1GDGD DRTR NORTH, NDRTH, SOIJTH, SOIJ
TH, WEST, WEST, ERST.، EFST
IOG10 DATR NDRTH ERST,NORTH ERST
,SOIJTH ERST, &OUTH ERST, SDIJTH WES
T, SOIJTH UEST, NDRTH WEST,NORTH WE
ST
```

Now that is starting to look rather silly as both arrays now contain exactly the same words, so why not get rid of the response array, R\$, and simply print $\mathrm{C} \$(\mathrm{~N})$ ? Well, in this case you could do that without any problem, but of course where the responses are not simply a repetition of the input (as is very often the case) the second array is essential.

If you look hard at all those subroutines you will realise that they all do only one thing - update the values of X and Y . Now we could include that information in the original DATA and get rid of them altogether! We need to add two more arrays to hold the $X$ and $Y$ coordinates, add the appropriate values into the DATA lines after eachresponse, and READ in this information in blocks of four (INPUT, RESPONSE, X-MOVE, Y-MOVE-see Table 2.2).

[^0]
## 10010 DRTR NDRTH ERST,NDPTH ERST

 , 1, -1, SDIJTH ERST, SDIJTH EAST, 1, 1 , SOIJTH WEST, SOIJTH WEST, $-1,1$, NORTH WEST, NDRTH WEST, $-1,-1$11016 REFD C

| ELEMENT <br> NUMBER | COMMAND <br> WORD <br> CS(n) | RESPONSE <br> RS(n) | X-MOVE <br> X(n) | Y-MOVE <br> Y(n) |
| :---: | :---: | :---: | :---: | :---: |
| 1 | NORTH | NORTH | 0 | -1 |
| 2 | SOUTH | SOUTH | 0 | 1 |
| 3 | WEST | WEST | -1 | 0 |
| 4 | EAST | EAST | 1 | 0 |
| 6 | NORTH-EAST | NORTH-EAST | 1 | -1 |
| 7 | SOUTH-EAST | SOUTH-EAST | 1 | 1 |
| 8 | NOUTH-WEST | SOUTH-WEST | -1 | 1 |

Table 2.2 X and Y Moves Incorporated into Arrays

Now we can delete lines 300 to 2700 and modify line 210 so that $X$ and $Y$ are updated here (see Flowchart 2.5).

##  

This overall pattern of putting all the information into a series of linked arrays is a very common feature which is used in several of the later programs in this book.


Flowchart 2.5 Using linked Arrays

## Abbreviated commands

So far we have always used complete words as commands, but that means that you have to do a lot of typing to give the machine your instructions. If you are feeling lazy y ou might think of changing the command words to the first letter of the words only, and then INPUT a single letter. However, unless you start using random letters that will only work as long as no two words start with the same letter! To code all the eight compass directions used above, we will have to use up to two letters: N, NE, E, SE, S, SW, W, NW.

D, 1, W, WEST, -1, (O, E, ERST, 1, B
1 DV1] DRTR NE, NORTH ERST, $1,-1$, SE
, SOUTH ERST, 1, 1, SW, SOUTH WEST, -1
1 , NW, MORTH WE8T, $-1,-1$

Notice that it is only the actual command words which have changed and that the computer gives a full description of the direction, as we are still using that second array which holds the response.

## Partial matching

In all the programs above we have always checked that the input matched a word in the command array exactly. However, it would be useful if we could allow a number of similar words to be acceptable as meaning the same thing. For example, you could check whether the first letter of the input word matched the abbreviated keyword by only comparing the first character (taking LEFT\$(IN\$,1)).

## 190 IN\%=LEFT(IN(1)

That will work with NORTH. SOUTH. EAST and WEST, but there are obvious problems in dealing with the intermediate positions. In addition there are lots of words beginning with the letters N, S, E and W - all of which would be equally acceptable to the machine as a valid direction.

For example:

## NOT NORTH

would produce:

## GOING NORTH

A more selective process is to match a number of letters instead of just one. In this example the first three letters of the four main directions are quite characteristic.

NOR
SOU
EAS
WES

If you use these as command words, then, for example:

```
        NOR
        NORTH
        NORTHERN
and NORTHERLY
```

will all be equally acceptable, but:

```
    NOT
    NEARLY
    NOWHERE
and NONSENSE
```

will all be rejected.
All we need to do is to take the first three letters of the input, LEFT\$(1N\$,3), and compare them with a revised DATA list. Line 10010 can be deleted and the word number variable WD must then be amended to 4 .

```
29 L[;-3
190 1N%wLEFT&(IN%,3)
10שOG ORTR NOR.NORTH.D, -1, SOU,SD
IJTH, D, 1.WES,WEST, -1, D,ERS,ERST,1
    .g
```


## Sequential commands

In the routines above we have dealt with the intermediate compass positions as separateentities, but if we could give a sequence of commands at the same time we would not need to do this. There is always more than one way to get to any point, and if more than one command word could be understood at the same time we would not have to worry about checking for directions such as 'NORTH EAST' as they could be dealt with by the combination of 'NORTH' and 'EAST'.

This brings us to the very significant question of how to split an input into words. First you must ask yourself how you recognise that a series of characters make up a separate word. The answer, of course, is that you see a SPACE between them. Now if we look for spaces we can break the input into separate words which we can look at individually. The easiest way to look for spaces is with the INSTR command whichsearches the whole of a designated search string for a match with a second target string.

For example, line 130 will check whether the first character in IN\$ is a space. If it is not a space then it will automatically continue checking until the end of IN\$ is reached. If no space is found in the whole of IN\$ then SP will be zero. If a space is found then the value of SP will be the number of characters along IN\$ that the space is located (see Flowshart 2.6).


Flowchart 2.6 Locating the Position of a Space

Try this out with:

## NOR WES

SP 4

## NORTH WEST

SP 6

## NOR NOR WEST

SP 4

Notice that the length of the word is accounted for by SP but that only the first space is found. To find all the spaces we are going to have to work harder. First delete that temporary line 140.

Let's look at the input logically from the start (lefthand side). We will replace the LEFT\$(IN\$,3) with MID\$(IN\$,ST,3) so that we can look at any three-letter combination in the whole of IN\$. To make it moresensible we will call the result of this W\$ as it shows the position of a word. Tostart with we must set the search start position ST equal to one and add a space to the front of IN\$ so that the first word is also found (see Flowchart 2.7).

```
125 ST=1:IN$=" "+IN*
130 SP=INSTR(ST,IN由." ")
190 W%=MIDM( IN%,ST,3)
```



```
(N):X=Y+X(N):YzY+Y(N):GOTO IED
```

If you run this as it stands then you will still only find the first word as we have GOTO 100 on the end of line 210 . However simply sending the program back to the INSTR check in line 130 instead does not help either, as it will always start checking from the beginning of IN\$ and will always find the same first space. Once we have found this first space we need to


Flowchart 2.7 Searching for a Keyword
move the start position ST for the next search on to the character after that space, SP+I. When no more spaces can be found then the end of the input has been reached and we can GOTO 100 again.

```
140 IF SP>D THEN ST#SP+1:GOTO 19
\emptyset
1!50 GOTO 100
210 IF W%=Cक(NY THEN PRINT G%jR$
(N): Y=X,4Y(N)}Y=Y\Y(N):GOTO 130
```

Now typing.

## NORTH WEST

produces:
GOING NORTH

## GOING WEST

and even:

NOR NOR EAST
is decoded as:

## GOING NORTH

GOING NORTH
GOING EAST

It would be a lot neater if we deleted allthose redundant 'GOINGs' and put all the reported directions on the same line. We need to PRINT G\$ once, immediately before the INSTR check. Now each time we go through the loop comparing the current word with those stored, we PRINT R\$(N); if there is a match. As there is a semi-colonafter this, the words will be printed on the same line but we also need to add spaces between them. Finally we add a simple PRINT just before we go back for a new input, to move the cursor position on to the next line.

```
126 F'FINT G$;
145 FRINT
210 IF W$=C$(N) THEN FFIINT F゙$(N)
;" ";:X=X X X X N ) :Y=Y +Y(N) =GUTO 1ZO
```

Now:

## NORTH EASTERLY SOUTH WEST

sends you neatly round in circles:
GOING NORTH EAST SOUTH WEST

## CHAPTER 3

## Understanding Natural Language

So far we have only communicated with the computer in a very restricted way, as it has only been programmed to understand a very few words or letters and it only recognises these if they are entered in exactly the right way. For example, if you put a space before or after your command as you INPUT it then it will be rejected. This is because we are comparing whether the two strings match exactly.

On the other hand in the real world everyone uses what is known as 'natural' language which is a very sophisticated and extremely variable thing which only the human brain can cope with effectively. Even if we forget for the moment the difference between 'English' and 'American' or even regional dialects of either (can 'Ow bist old but' really mean 'How are you old friend'?) dealing with language has an inf inite number of problems.

Even the most sophisticated systems in the world cannot cope with everything. There is an old story which illustrates this point very well. The CIA developed a superb translation program which could instantly convert English into Russian and vice versa. In the hope of impressing the President they laid on a demonstration of its capabilities, in which it converted everything he said into Russian, spoke that, and then retranslated the Russian back into English. He was most impressed and was totally absorbed until one of his aides reminded him that he had forgotten that the First Lady was waiting for him outside. When he ruefully commented 'out of sight, out of mind' he was amazed to hear the machine come back with 'invisible maniac'!

## Dealing with sentences

Everyone knows that real language is made up of sentences, but what exactly do we mean by a sentence? Well, the most obvious way we recognise a sentence is that we see a full stop! However if we are going to be able to deal with sentences, we are going to have to think a lot harder than that.

The Oxford Dictionary definition includes 'a series of words in connected speech or writing, forming grammatically complete expression of single thought, and usually containing subject and predicate, and conveying statement, question, command or request' but also concedes that it is used loosely to mean 'part of writing or speech between two full stops'. Phew! Can somebody translate that into everyday English, please?

The intricacies and illogicalities of the English language are infamous so how can we expect a computer to cope?

Well, let's start by looking at some simple examples of sentences.

## I WANT.

consists of a subject I and a verb WANT

## I WANT BISCUITS.

also has an object BISCUITS

I WANT CHOCOLATE BISCUITS.
qualifies the object with an adjective CHOCOLATE

## I SOMETIMES WANT CHOCOLATE BISCUITS.

qualifies the verb with an adverb SOMETIMES.
The most important word in all the above examples was 'WANT' as it conveyed the main idea. The second example was more informative as it indicated that only one particular type of object, BISCUITS, was wanted. The addition of an adjective, CHOCOLATE, gave further information on the type of object wanted, but life became more uncertain again when the adverb SOMETIMES was included.

Now how could a computer program decode such sentences? The answer must be to find some logical structure in the sentence, so what 'rules' could we lay down for this example?
I) All started with a subject I and ended with a full stop.
2) The last word was always the object BISCUITS (unless there was no object and only two words).
3) If the word before the object was not the verb WANT it was an adjective CHOCOLATE.
4) If the word before the verb was not the subject I it was an adverb SOMETIMES.

Let's write a program in which we give the computer sentences and ask it to break them up into their component parts.

To start off, we need to give it a vocabulary of objects, adjectives and adverbs to work with. We will READ thesefromDATA and storethem in arrays $\mathrm{OB}, \mathrm{AJ}$ and AV , according to type.

```
10 GOSUB 10004
10000 [IM OR&(号), R.J*(5), F\/%< 2)
19999 REM OEJECTS
110FD [ORTA BISCUITS, EIINS.CFKE
11019 [JRTA COFFEE,TER, WGTER
11ज19 PEM FCOJECTI VES
110:25 [JRTA CHOCOLRTE,GINGER, JRIM
1103G [,ATR COL[, HOT, LUKEWRPI
110.3 REM RDVERBS
11540 DATR RLWRYS, OFTEN, SOIMETIME;
S
11100 FIJR N=币 TO 5
11110 PERC CJ&%(N:
1112ด NEKT N
11130 FOP N=\sqsubseteq TO S
11145 PERD F.J%(N.
11150 NEX'T N
11160 FOP.Na0 TO 2
11170 RERC; R\*(N)
111g0 NE:\T H
11190 RETIJRH
```

Now we need to break the sentence intowords（see Flowchart 3．1）．Once again we will do that with an INSTR search for spaces，and to make life easier we will add a space onto the end of IN\＄so that the format of the last word looks just like that of other words．

100 INPIJT ING
120 IN第＝INぁ＋＂＂
130 SP＝INSTRiST，IN果，＂＂）
190 GOTO 130
The end of the sentence has been reached when no more spaces can be found．

## 140 IF SP＝0 THEN 206

If a space is found then the section of IN\＄from ST（current search start） to SP－ST（current space－current start＝length of word）is cut out and stored in a word store array W\＄（ WC）．

```
150 W$(WC)=NID$(IN$,ST,SF-GT)
10010 DIM W吊(4)
```

To begin with $\mathrm{ST}=1$ so that the search starts at the first character in the


Flowchart 3．1 Cutting Out Words
input string．The word count variable WC is set to zero so that the first word found is stored in the zero element of the word store array．

## $11 日 S T=1: W C=$ の

The word count is incremented（so that the next element ofthe array W\＄is used next time）and a check made that there are not more than five words in the sentence．The start position for the next search is then set to one more than the position of the last space and the search is continued．

159 WC＝WC＋1
17G IF WC，
TOD LDHC＂：GOTD 1以
196 ST＝SP＋1

A test is now made to see whether there is a match between the key words in the sentence and the objects in the vocabulary array OB\$(N) (see Flowchart 3.2). Only words 2,3 and 4 are checked as these are the only possible


Flowchart 3.2 L.ooking for a Match
positions for the object in our restricted sentence format. Three different routines are jumped to according to the position of the matching word in the sentence. If no match is found a message is printed and a new input requested.

```
290 FOR N=0 TO 5
210 IF W% 2)=0B%(N) THEN E以¢
220 IF Win( 3)=01B(N(N) THEN EDQ
235 IF W'$(4)=CDP(N) THEN TGO
34 NEYT N
#5g PRINT "UQ.JECT NOT FOIJND"
?50 FOTD 150
```

If the object was found as word 3 then there was neither adjective nor adverb.

## 50D PRINT＂MO RO．JECTIVE OR ADVER $B^{\prime \prime}$ 510 COTO 1019

If the object was found as word 4 then there could have been either an adjective or an adverb in the sentence（see Flowchart 3．3）．

EBD PRINT＂EITHER RD．JECTIVE OR A ［VYERB＂


Flowchart 3．3 Adserb or Adjective

First we check for a match between the second word and the contents of the adverb array

61日 FOR $N=$ ず TO 2
620 IF WHe 1 ）$=$ सV禺 $(N)$ THEN 900 E3D NEXT N

If no match is found then we check the third word against the adjective list

```
640 FOF N==0 70 5
65O IF W$(2)=AJ$(N) THEN 1000
660 NEXT N
```

If a match is not found in either of these lists，then it would be useful to
indicate which word was not understood. The simplest answer is to check whether the second word was not the verb 'WANT', as in that case the second word must have been an adverb. On the other hand, if the second word was the verb then the third word must have been an adjective. Notice that the actual word which did not match is now included in the message.

```
E.7D IF WGC 1 <<>"WRNT" THEN PRINT
"RDVERE ":WG? 1);" NDT UNDERSTOOD
" EI.SE FRINT "RD.IECTIVE ":W$(2);
" NOT IJNDEPSTDO[""
6?% COTO 100
```

If a match is found in either test then a success message is printed. Note that these possibilities are exclusive and that in four words we can only have one or the other.

```
90D PRINT "RDVERB"
910 GOTD 100
1000 PRINT "RD.JECTIVE"
1010 COTO 100
```

Where both adverb and adjective are present we must check for both, and therefore a match in the first test also jumps on to the second test (see Flowchart 3.4).

```
TOD PRINT "RDVERE RND RD.JECTIVE"
710 FOP N*D TO 2
720 IF W&(1)=RV$(N) THEN 750
730 NENT N
```

If no match is found for the adverb, then this fact is reported: a flag AV is set to $l$ to indicate failure at this point before the adjective is checked.

```
740 PRINT "RDVEPE ";WMC1);" NOT
IJNDERSTDOD": : RV=1
750 FOR N=0 TD S
7ED IF W&(3)=F.Js<N: THEN g0D
770 NEXT N
```

If a successful match for the adjective is not found then the program loops back after a report.


```
OT INNDEPSTDOD"
7aด GOTO 1月0
```



Flowchart 3.4 Adverb and Adjective

If the adjective was found then a test is made that the adverb flag AV was not set before a match is reported. In any case, the flag is reset before the next input.

```
OBD IE RV=0 THEN PRINT "ROSECTIV
E FIND RDLPERE DK"
810 AV=0
3PQ GOTD 190
```


## What about punctuation?

As we have already said, you usually recognise the end of a sentence because it has a full stop, although when you type into a computer you usually forget all about such trivialities. But what will happen in the
program so far if some 'clever' user puts in the correct punctuation? If you think for a moment, you will realise that the computer will start complaining as it will no longer recognise the last word, as this will actually be read as the word plus the full stop.

We therefore need to check if the last character in the input string IN\$ is a full stop: this is simple as the ASCII code for this character is 46 . The best place to check seems to be immediately after the INPUT. If the code of the last character is 46 , then simply throw this character away and then continue as before.


Flowchart 3.5 Dealing with Punctuation
We will add this as a subroutine which is jumped to as soon as an input is made. Other punctuation marks may also appear at the end of the sentence, so we will read the last character as a variable LC which we will also use later. This is stored as a simple variable by taking the ASCII code of the last character in IN\$: using simple variables saves a lot of typing of string (\$) indicators (see Flowchart 3.5).

## 195 EOSJB 2000

2ดดの LC=RSC(PITHTS(IN世, 1) ?
2ด10 IF LE 246 THEN $21 F W$
SF9F PETIJPN
 ETLIPN

More useful sentence terminators are the question and exclamation marks which often indicate the context of the words．We candistinguish these in the same way by their ASCII codes and，forthe moment，we will just report their presence．

## 2020 IF LC＝33 THEN PRINT＂EXCLRIMA TION＇＂：GOTO 210D 2929 IF Le：$=63$ THEN PRINT＂OIJEST10 N＂：GOTO 2100

The normal INPUT command will not accept anything after a comma， which it reads as data terminator．However，LINE INPUT will accept any text including commas．

## 10D LINE INPUT INE

Commas may be useful in indicating different parts of a sentence，which could be examined as＇s ub－sentences＇in their own right．However，in simple cases they are best deleted and replaced by spaces before the sentence is broken into words（see Flowchart 3．6）．Note that this will only function totally correctly if there is no space after the comma，as any space following a replaced comma will be seen as a new word．


Hlowchart 3．6 Replacing Commas and Apostrophes
115 COSIJB 39ดМ
3000 CIteINSTRCST，IN\＃，＂，＂）
3月10 IF CM＝0 THEN ST＝1：RETURN

HT\＄（ IN⿱，LEEN（INक）－CM）
3D30 ST $=\mathrm{Cl} 1+1$
उค45 COTO 30日ต

If you add this line, you can see the punctuation being taken out of the string item.

## 3 SOS PRINT INs

Apostrophes can be dealt with in the same way, except that we do not replace them with a space but simply close up the words.

115 GOSUE $3000:$ GOSUE 3100

```
3100 AF=INSTFi(ST,IN$,":")
S110 IF AF=0 THEN ST=1:FETUFN
3120 IN$=LEFT$(IN$, AF-1)+FIIGHT$(
IN$,LEN(IN$)-AF')
3125 F'RIINT IN$
S1SO ST=AF+1
3140 GOTO S100
```


## A sliding search approach

Although the method of examining a sentence described above will work, it has the disadvantage that it requires the sentence to be entered in a particular, restricted format. For example, if you enter:

## I WANT HOT CAKES OFTEN

the computer will report:

## OBJECT NOT FOUND

as it mistakenly takes the last word OFTEN as the object.
On the other hand using a sliding search of the whole sentence for each key word, without first breaking the sentence down into words, has the advantage that it allows a completely free input format. In this approach we take the first key word and try to match it against the same number of letters in IN\$, starting at the first character. If this test fails then it is automatically repeated, starting from the second character, etc, until a match is found or the end of IN\$ is reached. For example, if IN\$ was 'I WANT CAKE' and the first key word was‘CAKE' the comparisons would be:

| Pass 1 | I WA |
| :--- | ---: |
| Pass 2 | WAN |
| Pass 3 | WANT |


| Pass 4 | ANT |
| :--- | :--- |
| Pass 5 | NT C |
| Pass 6 | T CA |
| Pass 7 | CAK |
| Pass 8 | CAKE (match found) |

Delete lines 105-1010 and add this line to check for the first object OB\$(0).

## 210 SP=INSTR(1,IN:, OBW(N)):IF SP \O THEN PRINT OBE(N):" ";

Each object can be compared in the same way by forming a loop. (Note that printing a semi-colon after OB\$(M) ensures that each word is printed on the same line.)

## 200 FDR N=0 705 220 NEMT N

Similar checks can be made for matching with words in the adverb and adjective arrays.

```
SOO FOF N=O TC 2
310 SF=INSTF(1,IN$,AV$(N)):IF SF'
>0 THEN FYTINT AV$ (N);"
32O NEXT N
4OO FUF N=O TC) 5
410 SF=INSTF(1,IN$,AJ$(N)):IF SF
O THEN FFINT AJ$ (N):"
420 NEXT N
1000 GOTO 100
```

To report what has been found, and so that we can use the words discovered later, we will store each in an array as it is detected. We already have a word store array W\$ but we will expand it to hold up to 20 words (which should be enough for even a very verbose sentence!).

```
10010 DIM W$(19)
```

If a match is found a temporary string $\mathrm{T} \$$ is set equal to the matched word, and a subroutine called at line I500, which puts the worddetected in thefirst array element ( see Flowchart 3.7).

210 SP=INSTR(1,INF.OBक(N)):IF SP
" 0 THEN PRINT OB\& (N);" ": T\$aOB

## (N): GOSUB 1500 <br> 



Fluwchart 3.7 Sliding Search

The word count WC is then incremented, so that the next word is put in the next element. before returning.

## $15 ? 9$ WCx $\mathrm{FLC}+1$ <br> 153D RETURN

Using a temporary string $\mathrm{T} \$$ in the actual subroutine means that we can also use it in the tests for adverbs and adjectives in exactly the same way.

```
310 SF=INSTF(1,IN$,AV$(N)):IF SF
O THEN FFRINT AV$(N);" ";:T$=AV$
(N):GOSLUE 150O)
410 SF=INSTF(1,IN$,AJ$(N)):IF SF'
s0 THEN FFFINT AJ$(N);" ";:T$=:AJ$
(N):GOSUE 15OO
```


## Partial matching

One advantage of the sliding search is that you can easily arrange to recognise a series of connected words by only looking for some key characters. This is obviously useful as it saves you having to put in both single and plural nouns such as BISCUIT and BISCUITS. If you amend the DATA in line 11000 as shown below than both will be recognised.

110゙GO CATA EISELIT EUH DAYE

However life is not that simple as using BUN rather than BUNS can produce some unexpected results. On the plus side it will detect BUN. BUNS, and BUNFIGHT but unfortunately BUNCH, BUNDLE, BUNG ALOW, BUNGLE, BUNK, BUNION, and BUNNY as well'


Flowchart 3.8 Checking That This is the Start of a Word

This problem is not restricted to prefixes as the computer will also not distinguish between HOT and SHOT. You could include a check that the character before the start of each match was a space (ie that this was the start of a word, see Flowchart 3.8). SP gives the current start-of word position so MID\$(IN\$,SP1,1) is the character before this.


## (1) THEH IF MID( $⿻$ (INक, SP-1, $1 \geqslant="$






:cosue zoog

Y THEN IF MI CHE IN\&, $\mathrm{SP}-1,1 \geqslant="$
 : GOSiB 2Gug

For this to function correctly on the first word. just add a space to the start of IN\$.

## 110 INक=" " + IN

In a similar way. you could use checks on the next letter after the match, or the length of the word, to restrict recognised words.

## Putting things in order

Although we have now detected all the words in the sentence, regardless of
their position or what else is present, they are found and stored in the order in which they appear in the DATA. This is because the comparison starts with the first item in the object array rather than the first word in the sentence. It would be useful if we could rearrange the word store array so that the words in it were in the order in which they appeared in the sentence.

To do this, we must keep a record of the sentence position of the word SP and word count WC, as each word is matched in a new word position array WP. This is a two-dimensional array with the sentence position kept in the first element, WP(WC,0), and the word count, WP(WC,1), in the second.

```
10020 [IM WP(19,1?
151@ WF(WC,0)=EP:WF(WC,1`=WC
```

The actual sorting subroutine which doesthe rearrangement is at line 4000. This must only be reached if a match is found.

```
44F) IF WC=g THEN 4'口
450 C2OSUE 4900
460 GOTO 19円
470 PRINT"NG MATCH FOIJND"
499 GOTO 190
```

The sort routine (see Flowchart 3.9) takes the sentence position of the first word found (first element in the first dimension WP( 0,0 )) and compares it with the sentence position of the second wordfound (second element in the first dimension WP $(0+1.0)$ ). If the position variable for the first word is of


Flowchart 3.9 Putting Words in Order
higher value than that for the second word then the first word found is farther along the sentence than the second word, and these therefore need to be swapped around. This will put the sentence-position pointers right but the word-count markers also need to be rearranged to the correct positions. This process is repeated until the word pointers are all in the correct order. Notice that the actual contents of the string array which holds the words are not altered but only the pointers (index) to them.

```
40以リ FOP N=0 TO WC--2
401: IF WP(N,G)WWP(ri+1,0) THEN N
ENT N GOTO 4040
```



```
    WP(N+1,O
4030 D=WP(N,1):WP(N,I)=WP(H+1,1)
    WP(N+1,1)=0.GOTD 400日
```

If the strings are now printed in revised word－count WC order，they will be as they were in the original sentence，which should make it easier to understand them．

```
4040 PRINT:FGP N=FJ TOU WC--1
405FJ PRINT W$CWP(N,1) );" ";
49EGF PESTT N:PRINT
```

All elements in the sentence position array $\mathrm{WP}(\mathrm{N}, 0)$ and the word count WC must be reset to 0 before the next input．
4076 FOP $N=0$ TO ..... 194080 WP（N，の）＝■
4096 NEXT N
$41 \mathrm{GJ} W \mathrm{WC=}$
411D PETIJRN

## CHAPTER 4 <br> Making Reply

## More sensible replies

We have considered at length how to decodesentences which are typed into the computer, but the replies it has produced so far have been very limited and rigid. Although many of the original words in a sentence are of ten used in a reply, in a real conversation we look at the subject of the sentence and modify this word according to the context of the reply.

For example the input:

## 1 NEED REST

might expect the confirmatory reply:
YOU NEED REST
and similarly:

## YOU NEED REST

should generate:
1 NEED REST

If you look at that situation logically. you will realise that for each input subject there is an equivalent output subject. and that we have simply chopped off the original subject and added the remainder of the sentence to the appropriate new subject.
' I ' is only a single character so we could check LEFT\$(IN\$, I). If this was 'I' then PRINT "YOU" could be added to the front of the remainder of the input, RIGHT\$(IN\$.LEN(IN\$)-1).

```
10 INPUT IN%
3@ IF LEFT&{IN&,1)="I" THEN PRIN
T "YOU"+RIGHT$(IN果,LEN(IN*)-1)
60 [,0TO 10
```

In the same way，the first three characters LEFT\＄（IN\＄，3）could be checked against＇YOU＇and replaced when necessary by＇ I ＇．

## 50 IF LEFT（INq， 3 ！$="$ YOU＂THEN FR 

If you try that out with a series of sentences，you will see that it works OK until you type something like：

YOU ARE TIRED
which comes back as the rather unintelligent：

## 1 ARE TIRED

We could get around this by checking forthe phrases＇I AM＇and＇YOU ARE＇as well as＇I＇and＇YOU＇on their own，but notice that you must test for these first and add GOTO 10 to the end of lines 20 and 40 to prevent a match also being found with＇I＇and＇YOU＇alone．
 FINT＂YOIJ ARE＂＋RIGHT\＆（IND，LENV IN 43－4）：GOTO 19
 H PFINT＂I RM＂＋RIGHTHECIN，LENCIN क ）－7）：GOTO 10

Although this method will work，the program soon gets very long－winded as a separate line is needed for each possibility as we must take into account the length of the matching word or phrase．Where many words are to be checked，it is therefore better to use a multidimensional string array which can be compared with the input by a loop．

A convenient format is to have a two－dimensional array $\operatorname{lO}(\mathrm{n}, \mathrm{m})$ where the first dimension of each element， $10 \$(n, 0)$ ，is the input word or phrase and the second dimension， $10 \$(n, I)$ ，is the corresponding output word or phrase．It is easier to avoid errors if these are entered into DATA in matching pairs and READ in turn into the array．Start a new program with these lines which set up the array．

10 GOSUE 19の日の
10以リの DIM IOథ（3，1）

ARE，YOIJ ARE，I FIM
129日g FOR N＝ด TO 3
12910 FERD ID\＆（ト，O）III世（H，1）
12 O 3 G NEST H
139日g FETIJRN


Flowchart 4.1 Using a Corresponding Reply

We will use a looping sliding string search again，which for the moment will just print out the corresponding word or phrase to that matched，IO \＄（N，I） （see Flowchart 4．1）．One advantage of the sliding string search here is that it will happily match embedded spaces in phrases as we have not broken IN\＄ into＇words＇before matching．

```
189] LINE INPUT IN&
Z09 FOP N=0 TO 3
21月5P=INSTR(1,IN昨,ID$(N,O)):IF
SP`ज口 THEN PRIMT IO&;N,1?
2%ด NEXT N
25ด [.OTO 100
```

It is better to redefine the required response word as a new string which is the first part of the reply R1\＄，and then PRINT this when the loop is left．

##  <br> SP＞D THEN PI $\$=10 \%(N, 1)$ <br> 230 PRINT R1＊

To get a fuller reply，we need to add back on the rest of the original sentence R2\＄（after inserting a space）．It is not difficult to define the＇rest of the sentence＇．We just need to subtract the end position of the word from the LENgth of the sentence and use this value in RIGHT\＄．SP points to the start of the matched word：we have a record of the LENgth of this word in the first dimension of the array as $1 O \$(N, 0)$ ，so we just need to subtract SP＋LEN（IO\＄（N，0））．




- 日) $3>$

20 FROUT RIt+F2


Flowchart 4.2 A Fuller Reply

Now when you try:

## I AM CLEVER

the computer agrees:
YOU ARE CLEVER

But if youthen press RETURN againit stilltells youthat you are clever which is not true, as you have not emptied R1\$ and R2\$ before looping back to the next input!

Before you feel too clever try:
WE ARE STUPID
which may well surprise you when it gives the reply:

YOU

If you think for a few moments, you will see that one of our keywords is hiding inside another word in this particular sentence. If you cannot see it then try:

## WE ARE INCOMPETENT

where the computer disagrees with you by returning:

## YOU COMPETENT

Although each keyword is tested for in turn, each one is set to RI\$ when a match is found so only the last match is reported. As the keyword is only checked for once in each sentence, embedded 'l' only causes problems when this is not the keyword in the sentence.

To get around this we must consider which keywords may cause problems. Although the letter ' I ' is very common, it is very rarely the last letter in a word and so we could check that there is a space after the keyword. We must treat all keywords in the same way so add a space to the end of each. This could be done by changing all the DATA but it saves memory in the long run if we add the space as the array is set up. Note that there is no need to add spaces on to the end of the replies.


We also now need to subtract one less character from IN\$ to give R1\$, as the space has now become an integral part of the keyword.

```
210 SF=INETR?1,IH*,ID$(N,O%):IF
SF>G THEN R1%=IO*(N,1)\P2$=" "+F
```



```
|(|)3)+1)
```

The computer will now readily agree about your incompetence.
If the first keyword is not at the start of the sentence, then everything before it will be ignored in the reply.

For example the answer to:

## WHAT IF I FALL?

will be:

## YOU FALL?

Some strange results can still occur when two true keywords are present. For example:

WHAT IF YOU AND I FALL
gives

I AND I FALL
and

WHAT IF I AND YOU FALL
replies

I FALL
However, adding more suitable keywords is easyand some combinations will just not be acceptable. To make the routine more general, it is better to define the number of keywords as a variable KW and use this in place of the actual number.

```
1F KW=5:GOSUE 1คю@ी
%W FOR N=0 TIJ YW
1GOF& [IM IO$(KW,1)
1151G [JATA WE,WE, THE'', THE''
120Н゙5 FIDP N=:の TO KM
```

Now the answer to:

WHAT IF WE FALL'
is the more logical:
WE FALL"

## Pointing to replies

So far our computer has displayed only slightly more intelligence than a parrot. as it has merely regurgitated a slightly modif ied version of the input. The next stage, therefore, is to make it take some logical decisions on the
basis of the input before it replies．First，we need to clear enough string space（ 1,000 bytes）and then jump to an initialisation routine．

## 10 CLERF 10105 <br> 2J GOSIE：10円ค月

The numbers of subjects SU，verbs VB and replies RP are defined as variables so that the program can be easily expanded，and three arrays using these are set up．（As we have a zero element in the array，these values are all one less than the number of words．） $\operatorname{SU} \$(n, n)$ is a two－dimensional array which is concerned with the subjects in the input and output sentences．The first dimension（ $\mathrm{n}, 0$ ）contains the recognised subject words and phrases allowed in the input，and the second dimension（ $\mathrm{n}, \mathrm{l}$ ）contains the opposites which may be needed in the output．VB\＄（n）holds the legal verbs，and RP\＄（n）a series of corresponding replies．

10015［IM SUq（S1．J，1）
19の？ด DIM vB\＆UE ？
1月以30 DIM RP末（RP）

Table 4．1：Pairs of Subjects in $\operatorname{SCS}(n, n)$

| SU\＄（n，0） | SU\＄（n，l） |
| :--- | :--- |
| I HAVE | YOU HAVE |
| I＇VE | YOU＇VE |
| I AM | YOU ARE |
| I＇M | YOU＇RE |
| YOU HAVE | I HAVE |
| YOU＇VE | I＇VE |
| YOU ARE | l AM |
| YOU＇RE | I＇M |
| YOU | I |
| SHE HAS | SHE HAS |
| SHE IS | SHE IS |
| SHE＇S | SHE＇S |
| SHE | SHE |
| THEY＇VE | THEY＇VE |
| THEY ARE | THEYARE |
| THEY＇RE | THEY＇RE |
| THEY | THEY |
| HE HAS | HE HAS |


| HE IS | HE IS |
| :--- | :--- |
| HE'S | HE'S |
| HE | HE |
| WE HAVE | WE HAVE |
| WE'VE | WE'VE |
| WE ARE | WE ARE |
| WE'RE | WE'RE |
| WE | WE |
| I | YOU |

The first two lines of DATA contain paired input a nd output subjects (see Table 4.1) and these are READ into corresponding dimensioned elements in the $\operatorname{SU} \$(n, n)$ array. As the pronouns ('I', 'YOU'. etc) are frequently linked to other words to form phrases (such as 'l'VE'), these combined forms are also included in the DATA. Notice that these are arranged in such an order that the most complete phrase containing a keyword is always found first. A space is added on to the end of each element, so that some clashing of partial matches is avoided and a space is automatically formed in the reply.

11 GGa DATA I HAVE YOUS HANE, I'ME, 'rOU'VE, I RM, YOU RRE, I'M, 'OII'RE, Y OUS HA'VE, I HAVE, ‘OIJ" VE. I "VE, YOU R FE, I FM, YOIJ"RE, I M, MOIJ, I
11010 [JATA SHE HAS, SHE HAS. SHE I S, SHE IS, SHE'S,SHE'S,SHE, SHE, THE Y'VE. THE's"VE, THEY RRE, THEY RPE. T. HE'Y'RE, THE'r'RE, THEY, THEY
1 1mip lafta He HAS, HE HAS.HE IS.H F IS.HE'S.HE'S,HE,HE, WE HAVE, WE HAVE. WE' YE, WE'VE WE RRE, WE RFE. W E'RE, WE'PE, WE WE I I, YOIJ



, $19=\mathrm{St}$ 出 $(N, 1)+" \quad "$
1293日 HEXT N

The next DATA line contains the main verbs which are READ into $\mathrm{VB} \$(\mathrm{n})$. The verb 'to be' is omitted as its variations are so complicated. and many of its versions are already accounted for in the 'subject' check.

11530 DATA HATE, LDVE,KILL, DISLIK E,LIKE,FEEL, KNGJ

## 12041 FUF Nabj TO VE $12 ด 50$ PERD VE\& $N$ ) <br> 12565 NEXT N

The last set of DATA contains the replies which are put into RP\$(n), before control RETURNs to the main part of the program. To makethingssimple to understand and check at this stage, all the replies contain the original verb, although of course they could say anything.

```
11040 DiATR PFIIRREL'i' HRTE 'rOIS AS
WELL,L.OU/E YOU TOO,KILL. YOIJ,[MSLI
KE LOTS OF THINHS,LIKE CHIPS.FEE
L FDWNEPFIJL, KNOIN EUEF:`THING
120PG FOR N=0 TO RF
12\OmegaSO RER[ PP*(H)
12gGF NEMT N
13FDG RETIJRN
```


## Stripping the input

The input string IN $\$$ is searched for question marks, exclamation marks and apostrophes ( $\operatorname{CHR} \$(34))$, and these are cut out as before by a subroutine at line 1000 .


```
    "
110 ST=1:S尔">":GOSIJE 10@A
120 ST=1:S4=" "": TOS|E 10DE
```




```
101G IF SP=Dि THEN RETIIFH
```




```
103¢ ST=8P+1
10.4F GOTD 10@FI
```


## Matching

The input string is now compared with the list of subjects in the first dimension of $S U \$(n, n)$ (see Flowchart 4.3). If there is no match then a new input is requested, or else a subject match variable SM is set to the element number at which a match was found.


Flowchart 4．3 Setting Match Pointers

210 SF＝IHSTFR 1，INHi，SIJ（N，G））
2？0 IF SP＝日 THEN NEMT N：GOTO 195
$230 \mathrm{SM}=\mathrm{N}$

The verb array is now compared with IN\＄．If no verb is found，then the input is rejected，or else the verb match variable VM is set．

240 FDP $N=\square$ TO VE

2EG IF SF＝曰 THEH NE＇KT H：WTO 19日 270 $1 \mathrm{M}=\mathrm{N}$

## Making reply

Now that the subject and verb have been identified, we can pick up the appropriate reply by using $V M$ as a pointer to the reply array RP ( $n$ ).

## 

In the simplest case we can just add the appropriate subject to the front of RL\$ before we print it.

TrD PFIIT PL. $\$$
ᄃIF

Now, for example, if you type in:

## 1 HATE COMPUTERS

the program will reply with:

I PROBABLY HATE YOU AS WELL
and:

I KNOW A LOT
generates:

I KNOW EVERYTHING

## Alternative subjects

If you prefer the machine to agree withyourather thantrying to beat youat your own game, then just change the subject added to RL\$ to the second element of the array (the 'opposite').

now

I KNOW A LOT
generates:

YOU KNOW EVERYTHING

For more variety, you can pick the subject at random from the first or second element, so that the reply is not predictable.

```
510 PS=PHOC? -1
```



## Putting the subject in context

It would be more sensible altogether if we chose the correct subject according to the context of the reply, but to do that we must have markers in the reply array. We will use a slash sign '/' to indicate that the word in the first dimension of the subject array is to be used, and an asterisk '*' to indicate that the word in the second dimension is to be used.

```
1104E [ORTA PROEREL'' HRTE YOLI FS
```



```
ISLIKE LOTS DF THTHTG,, LIKE CHIP
S. FFEEL FRWWERFUL., *YHOW EVER`THIN
[1
```

We can now search the reply string $\mathrm{RP}(\mathrm{VM})$ pointed to by the verb marker VM for a slash sign ' $\%$ '. If one is found, then the contents of the first dimension of the subject array SU\$(SM.0) are added to the reply RL\$, less the first character (the slash sign. see Flowchart 4.4).


52円 IF SP\% THEH ODD

(FL(.) >-1)
S1ヵ . 10 OTO 530


Flowchart 4.4 Putting the Subject in Context

On the other hand if no slash sign is found in the reply then a second search is made for an asterisk＇＊＇．If this is found，then the second dimension of $\operatorname{SU} \$(n, n)$ is used in the same way．

```
590 SP=INSTP(1,FL彷,"*")
540 IF SF>日 THEN 820
```



```
(PL&)-1)
8.30 F.20T0 5.50
```

Now：

I LOVE ME
will give：

## I LOVE YOU TOO

but：

I FEEL POWERFUL
produces：
YOU FEEL POWERFUL

## Inserting into sentences

To make things simple，we have always started our reply sentences with the subject，but in real life this is not always the case．Now that we have markers in the replies to indicate what type of subject is to be added，we can also use them to indicate where in the reply to insert this word or phrase．First we will amend the DATA so that the word to beinserted is never at the start，to make the insertion process obvious．

```
1104日 [,ATA [GO YOIJ FEALISE THAT
PFOEFEL'Y' HATE `'OIJ FS WELL. WELL
L.DVE YO!J TOO, IF 人[ON|'T KILL 'MOו|
FIPST,SIJ UHAT *[ISL.IYE LOTS DF T
HINCSS,DO LIKE CHIPS,WHW COU :FEE
LPOWERFUL,HOW DO WFHOW EVEFYYTHIN
C-
```



Flowchart 4.5 Inserting into a Sentence

We actually already have a record of where to insert the word as SP tells us where in the reply the slash or asterisk was found. All we need to do is to take the part of the reply before the marker, LEFT\$(RL\$,SP I), add the correct version of $S U \$(S M, n)$, and then the rest of the reply RIGHT\$(RL\$,LEN(RL\$)-SP).

##  ) + RIGHT\$(PL. , LEM RL ( ) -SP)

##  

Now:

I WILL KILL HIM
produces:
IF I DON'T KILL YOU FIRST
and:
I DISLIKE COMPUTERS
gives:

Although we are now inserting the subject into the reply sentence more naturally, we are only dealing with one subject per sentence. Another minor modification will allow us to insert any number of subjects into a sentence. All we need to do is to keep repeating the search for markers until no more are found. A start variable ST is defined as I in line 500 , and then a search is made for the first type of marker. When a match is found, ST is reset to one more than the match position. When $\mathrm{RL} \$$ has been modified by line 800 we now need to jump back to 510 to look for more markers. If no match is found for the first marker then ST is reset to 1 . The secondty pe of marker is then checked for in the same way.

```
50ด RL."=PP&(UM1):ST=1
519 SP=INSTR<ST,RL$,"ノ">
52G IF SP`| THEN ST=:SP+1:GOTO 80
0:ELSE ST=1
```



```
540 IF SP>G THEN ST=:SP+1/GOTO O2
\square
910 GOTO 51@
830 GOTO 5.3n
11040 [ATA DO 'YOIJ RERLISE THRT,
FPDEREL.Y HRTE YOIJ RS WELL,WELL,
L.OQE YOU TOD, IF /DON'T KILL 'IOI
FIPST,SO WHलT /DISLIKE LOTS OF T
HINGS ESPECIRLLYY (DO LIKE CHIP
S,WHY [O *FEEL PDWERFIJL,*THINK *
F'NOLS EVERYTHING
```

Now:

I KNOW EVERYTHING
produces:

YOU THINK YOU KNOW EVERYTHING
and:

I DISLIKE COMPUTERS
gives:

## OBJECTions on the SUBJECT

Everything is starting to look rosy until you try something like：

## I HATE YOU

which replies：

DO YOU REALISE THAT YOU PROBABLY HATE YOU AS WELL

The problem here is that we are jumping out of the search routine as soonas the first match is found，and that although we are checking for the subject ＇I＇we are finding the object＇YOU＇first．As＇YOU＇comes before＇l＇in the subject array this is found first，in spite of the fact that it comes later in the sentence．

As we cannot practically mimic all the intricacies oft he human brain，we will have to make the assumption that the subject always comes before the verb，and the object after it．In the program so far we have been checking for the subject before we checked for the verb，and we will have to reverse that order．The verb position in the input is the value of $S P$ when a verb is found，so we will save that as a verb position VP pointer．

```
2の@ FOR N=@ TO WB
210 SP=INSTR(1,IN%, VR%(N)?)
2E!G IF SP=0 THEN NEYT N:GOTD 10口
D9g VM=N:VP=SP
```

Now when a match with the subject array is found，we can compare that position SP with the stored verb pointer VP，and reject the match if the subject is positioned after the verb（see Flowchart 4．6）．

```
340 FOF N-M TO SU
(5) SP=INSTP(1,INW,SU*(14, D`)
#E卬 IF SP=EF THEH NE&T N:COTO 19G
2TG IF SP%UP THEN NEMT N: FOTO 10
@
200 5M=14
```

（If you are too lazy to retype those lines you can add a couple of jumps to rearrange the order instead．）

140 โTOTO 249
231 GOTO 5月9
271 поTO ？ 29
2T9 UIM－N•UP－5P

## 235 IF SP:YP THEN NEXT H:GOTO ID ด



Flowchart 4.6 Rejecting Object Matches

## A change of tense

If we change to the past tense of the verb, we may or may not find this. With the first five verbs the situation is straightforward: to change to the past tense we just add ' $D$ ' to the end of the present tense. Both forms are ther efore accepted.

| HATE | HATED |
| :--- | :--- |
| LOVE | LOVED |
| KILL | KILLED |
| DISLIKE | DISLIKED |
| LIKE | LIKED |

However, with the last two ver bs the word changes completely, so there can be no simple match. Although we might get away with checking for ' KN ', as this is a rare combination, it would not be practical for us to use such a common group as 'FE' as a keyword.

| FEEL | FELT |
| :--- | :--- |
| KNOW | KNEW |

It is easier if we treat all verbs in the same way and, if there are no constraints on memory, then we can simply put all the possible versions into the verb array in pairs.

```
10Н0Н SU=2E:VE=13:PP=6
11日3D DATR HATE, HATED,LOVE,LDVED
,YILL.KILILED,DISLIKE,DISLIKED,L
IKE,I.IKED,FEEL,FELT, KNOIW, KNEW
```

Untess we want to have different replies for the different tenses, we will now have to divide the verb variable VM by two, to point to the correct reply for both forms.

## $237 \mathrm{WM} N / 2: W P=5 F$

## CHAPTER 5 Expert systems


#### Abstract

A human expert is someone who knows a great dealabout a particular subject and who can give you sensible advice ('expert opinion') on it. Such expertise is only acquired after long training and a great deal of experience, so unfortunately real experts are few and far between. In addition they are often not on hand when a problem needs to be solved.

Scientists have therefore a pplied themselves to the problem of producing computer programs which mimic the functions of such human experts. Such programs have the advantage that they can be copied very easily to produce an infinite number of experts. and of course they do not need teabreaks. sleep. pay-rises, etc. either! Of course. the computer must be totally logical and can still only follow pre-programmed instructions entered by the programmer. It is interesting to note that science fictionauthors have envisaged problems when the ultimate 'experts' (such as HAL in '2001: A Space Odyssey’ or Isaac Asimov's positronic robots) are faced with alternative courses which conflict with more than one of their prime directives and which produce not system crashes but 'pseudo-nervous breakdowns'.

Before we can start writing programs for 'expert systems'. we must ask ourselves how a human expert works.

Let us first consider the simplest situation, where the expert's task is to find the answer to a known problem.


First of all he takes in information on the current task.

Secondly he compares this with information stored in his brain and looks for a match.

Finally he reports whether or not a match has been found.

What we need here is simply a database program which tries to match the input against stored information (see Flowchart 5.1). A user-friendly system would accept natural language (see earlier). but to keep things simple here we will stick to a fixed input format. To start with. let's look at recognising animals by the sound they make. We set up two arrays: the question array $\mathrm{QU} \$(\mathrm{n})$ contains the sounds which are known. and each


Flowchart 5．1 A Simple＇Expert＇
element of the answer array $\mathrm{ANS}(\mathrm{n})$ contains the name of the relevant animal．


15G1G［OFTH MIFOW，OFT，WIIFF，［MG，MO
I．E＇IW，HOOT ，DWL，ME IT，H，HIN： HE
10，

1BW2の FETURY

Now we just need to ask for a sound and compare it with the contents of QUS（n）．

20）PFINT＂WHAT HOISE DOES IT MARE II．


EN 150
EF HE：\＆T H
EG PR［HT＂SOPF＇G I［OIH＇T KHOW THAT 마拒＂
ぞリ にロTロ ？

N；＂S IS A＂：Filto N）
110 ETOTO

Perhaps we should say at this point that our computer expert may well be better at this task than the human kind, as it cannot make subjective judgements, become bored, or accidentally forget to check all of the information in its memory. On the other hand it is not very literate as it reports 'A OWL', etc. (We will leave you to tidy that up by adding a routine which checks whetherthe first letter of the answerarray match is a vowel.)

## Branching out

The example above is very simple as only onequestionisasked, and there is only one possible answer. In reality we need to be able to deal with more difficult problems, where the answer cannot be found without asking a whole series of questions. For example, what should an expert do if he put the key in the ignition switch of his car and turned it, but nothing happened?
There could be a number of reasons for this:

FLAT BATTERY<br>BAD CONNECTIONS<br>SWITCH BROKEN<br>STARTER JAMMED<br>STARTER BROKEN<br>SOLENOID BROKEN

To find the cause. he should follow a logical path and make a number of checks. The first thing to do is to check whether it is only the starter motor which is not working:

## IS IGNITION LIGHT ON? (Y/N)

If the answer to this is ' $N$ ' then there is no power at the switch. so the cause must be one of the first three possibilities listed above. We can narrow things down more by finding out if the lights work:

## DO L.IGHTS WORK CORRECTLY? (Y/N)

If the answer is yes, thenthe battery cannot beflat and it must be connected to the light switch correctly. So presumably the switch is broken and a suggestion can be made that you replace it.

## REPLACE IGNITION SWITCH

If the lights do not work, then the connections should be checked.

## ARE BATTERY CONNECTIONS OK? (Y/N)

If the answer is yes. then the battery isflatso you must charge it (or push!)

## CHARGE BATTERY OR PUSH CAR

In the same way, a sequence of checks could be made to deal with a situation where there is power but the starter mechanism itself does not work (the last three possibilities).


Flowchart 5.2 A Branching 'Expert'
The simplest way to program this branching structure is by a series of IFTHEN tests (see Flowchart 5.2).

## 10 PFIMT"FFFILLT DIEATHISIS" <br> 2曰 PFINT

```
30 FRINT"IS IGNITION LIGHT ON& `'Y
小人"
40 INFIJT I&,螅
5@ IF IN的="Y" THEH ISQ
GE FRINT"DU LIGHTS WOPK COPPECTL
Y('゙州`"
70 INFUT IN$
36 IF IN年="Y" THEN 119
g| PRINT"REFLACE IGNITION SWITCH
1%G RUN
110 PRINT"ARE BATTEFY COHNEETIOH
50r (Y/N""
1%0 INPUTT IN垧
130 IF IN"*="Y" THEN 1EQ
149 FFINT "REFAIR EEtwIECTIONS"
15g FUH
160 FRINT"CHARGE EATTERY OR FUIEH
ERRI"
17日 PUN
180 --.-- etr --...--m
```

This sort of program is relatively easy to write．but as usual is inefficient as it becomes longer and more complicated．

## Pointing the way

A more efficient way to deal with the situation is to put the text into arrays and have pointers which direct you to the next question or reply，according to whether you answer yes or no to the current question（see Flowchart 5．3）．

The format for entering the DATA for each branch point is，then：
（TEXT），（Pointer for＇YES＇）．（Pointer for＇NO＇）
The first question was：
IS IGNITION LIGHT ON？（Y／N）．．． 1
If the answer was＇$N$＇then you need to ask the second question：
DO LIGHTS WORK CORRECTLY？（Y／N）．．． 2
Otherwise you need to continue with the other part of the diagnosis（which


Flowchart 5．3 Pointing to the Vext Output
we have not included but which would be point 7）．
We need to set up three arrays：OP\＄（n）contains the output（text）．Y（n） the pointer for＇yes＇，and $\mathrm{N}(\mathrm{n})$ the pointer for＇no＇．To make the program easy to modify，a variable NP is used for the number of points．The DATA is read in groups of three into each element in these arrays．Where the DATA point is a possible end of the program．this is indicated by the $Y(n)$ and $\mathrm{N}(\mathrm{n})$ pointers being set at zero．

1006Ct $\mathrm{HF}=7$

11 OLDD［OATA＂IS I INITIDH L．IGHT Dr
＂，7，2
11以10 DFTH＂ON L．I SHTS WORK SDPFE
にTL＇チ＂，3，4
1192日［JFTA＂PEFLALE SWITEH＂，ロ，ロ
 ロ゙tS OK＂，J，E

```
11g4B DATA "CHAFIE ERTTEFY DIR FIJ
SH EAR",O,G
11050 DAFTA "REPAIR EONNESTION",G
G
1105a DRTA "-rest of proarmm-",0
    0
12OFG FDR N=1 TO NP
```



```
12ด%ด NEXT N
13OFE FETIJFN
```

The actual runningroutine is very simple. Apointer CP is used to indicate the current position in the array: to begin with this is set to 1 and the first text printed. If this is an end point $Y(C P)=0$ (hardly likely just yet!), then CP is reset to I so that the sequence starts again. On the other hand, if a real pointer is present then an INPUT is requested. If the input is ' Y '. then CP is set to the value contained in the appropriate element of the $Y(n)$ array. otherwise it is set to the value contained in the $N(n)$ array.

```
20CF=1
3GJ PRINT DP程CP)
40 IF 'rCP%=G THEN 20
```




```
PGCF=F&CF
90%010 20
```


## A parallel approach

An alternative to the sequential branching method described above is the parallel approach which always asks all the possible questions before it reaches its conclusion. This method usually takes longer than following an efficient tree st ructure, but it is more likely to produce the correctanswer as no points of comparison are omitted.

Let us consider how we might distinguish between various forms of transport.

We will consider eight features and mark 1 or 0 for the presence or absence of these in each of our five modes of transport (see Table 5.1). If you look closely you will notice that the pattern of results varies for each of the different possibilities so it must be possible to distinguish between them by these features.

Table 5.1: Presence or Absence of Features

|  | bicycle | car | train | plane | horse |
| :--- | :---: | :---: | :---: | :---: | :---: |
| wheels | 1 | 1 | 1 | 1 | 0 |
| wings | 0 | 0 | 0 | 1 | 0 |
| engine | 0 | 1 | 1 | 1 | 0 |
| tyres | 1 | 1 | 0 | 1 | 0 |
| rails | 0 | 0 | 1 | 0 | 0 |
| windows | 0 | 1 | 1 | 1 | 0 |
| chain | 1 | 0 | 0 | 0 | 0 |
| steering | 1 | 1 | 0 | 1 | 1 |

We will enter these values as DATA and then READ them into a twodimensional array $\operatorname{FE}(n, n)$ which will hold a copy of this pattern, together with a string array containing the names of the objects $O B \$(n)$.

```
19 GOSUP 1EMAD
1000% [IM DE*(5),FE(5.8)
11000 DRTR EICYCLE,1,0,日,1,0,0,1
,1
11010 [,RTA CRP,1,0,1,1,0,1,0,1
110?g DATR TRAIN,1,0,1,0,1,1,0,0
11030 DRTA PL.ANE,1,1,1,1, [1,1,0,1
11040 DRTA HOPSE,D,日,0,0,0,0,0,1
1200D FOR N=1 TO S
12910 RERC OB$<N:
12ด20 FOF }M=1\mathrm{ TO &
12930 RERD FE<N,M`
12040 NEMT M,N
130gD RETIJPN
```

We can now ask whet her the first feature is present or not and use the reply to print out which modes of transport match at this particular point (see Flowchart 5.4).


```
SFG IPMPIJT ING
```



```
5%\ FOR N=1 T0,5
530 IF FE(N,1)=FNH THEN FRINT OR&
CN)
T4F NE:TT N
```



Flowchart 5.4 A Parallel Approach

In this case, answering ' $\gamma$ ' will produce a print-out of

## BICYLE

CAR
TRAIN
PLANE
and answering ' N ' will produce a print-out of only:

## HORSE

This clearly demonstrates a possible disadvantage of the parallel method as, although we have just shown that only a horse does not have wheels, the program insists that we still ask all the other questions before it commits itself. This is not really as silly as it seems at first, as if you answer ' $Y$ ' to the next question ('does it have wings') you will see that the computer quite logically refuses to believe in flying horses.

If we put the actual comparison part as a subroutine we can use it to check for all eight features in turn. We would need to make slight modifications, adding an array pointer AP which is incremented to check the next element of the feature array $\mathrm{FE}(\mathrm{N}, \mathrm{AP}$ ) in each cycle (see Flowchart 5.5).


Flowchart 5.5 Checking the Features in Turn

```
10日 PFIHT"DIDES IT HA'/E WHEELS"
119 GOSJP 50g
120 PFINT"DDES IT HRVE WINES"
130 [TJSNJP 50¢
145 PRIMT"DDES IT HF'VE FH& ENGIME
"
5% cosuse 509
150 FFIF{T"DOES IT HAVE TYRES"
179 TOSIJR 5AD
1QN PRINT"DDES IT HEED RRILS"
190 TOSUE 50G
```

```
ご以@ PPINT"[GOES IT HRVE WINCOWS"
Z16 TOGIJE 500
2?@ PRINT"DOES IT HR'v'E ff CHRIF{"
230 G0Sue 50ロ
246j PRINT"IS IT STEEFRELE"
#以 TOSIJE 50G
4DIJ PPINT
4 1 0 ~ R U N
515 RP=RP+1:RN=1:IF It拈*"N" THEN
    *1N=0
530 IF FE(N,AP)=RN THEN PRINT OE
$(N)
5ED PETUPN
```


## Top of the pops

The previous routine will print out a list of matches for each individual question as it proceeds，but it will not actually tell us which set of DATA is an overall match for the answers to all the questions．We can produce a SCORE which shows how well the answers match the DATA by having a success array element $\mathrm{SU}(\mathrm{n})$ for each object，which is only incremented when a match is found $\mathrm{FE}(\mathrm{N}, \mathrm{AP})=\mathrm{AN}$（see Flowchart 5．6）．


Flowehart 5．6 Measuring Success

```
260 PRINT
279 PRI&\"SCOKE"
2 8 9 ~ P R I N T ~
30G FOR N=1 TO 5
```

```
31ด PRINT DP每(N).SU(N)
320 NEMT N
530 IF FE(N,AP \=FNN THEN PRINT DE
$(N):SU(N)=SU(H)+1
15010 OIM SIJ(5)
```

If a complete match is found then $\mathrm{SU}(\mathrm{n})$ will be equal to 8 . Where one or more points were incorrect the score will be lowered. Scoring in this way is particularly useful where the correct answers to the questions are more a matter of opinion than fact (eg is a horse really steerable?), as the highest score actually obtained probably points to the correct answer anyway. (Notice that in this case each correct answer has equal weighting.)

## Better in bits

You may have noticed that we just happened to use eight features for comparison and it may have occurred to you that this choice was not entirely accidental as there are eight bits in a byte. If we consider each feature as representing a binary digit (see Table 5.2), rather than an absolute value, then each object can be described by a single decimal number which is the sum of the binary digits, instead of by eight separate values. We will convert to decimal with the least significant bit at the top so that, starting from the top at 'wheels', each feature is equivalent to I, 2, 4. 8, 16. $32,64,128$ in decimal notation.

Table 5.2: Binary Weighted Features

| bicycle | car | train | plane | horse |
| :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 1 | 1 | 0 |
| 0 | 0 | 0 | 2 | 0 |
| 0 | 4 | 4 | 4 | 0 |
| 8 | 8 | 0 | 8 | 0 |
| 0 | 0 | 16 | 0 | 0 |
| 0 | 32 | 32 | 32 | 0 |
| 64 | 0 | 0 | 0 | 0 |
| 128 | 128 | 0 | 128 | 128 |
| 201 | 173 | 53 | 175 | 128 |

It is not too difficult to convert our 'score' of I to 8 into the appropriate binary value, as long as we remember that the decimal value of the binary digit BV must double each time we move down, and that we must only add the current binary value to the score if the answer was 'yes' ( $A N=1$, see Flowehart 5.7).


Flowchart 5.7 Producing a Binary Score
If you consider for a moment, you will realise that we only need to keep track of the total number produced, SU , by adding the binary values of the 'yes' answers. There is no need to loop through and check each part of the array contents each time, or even to have a two-dimensional array at all! The only DATA we need to enter are the overall decimal values for each object, $\operatorname{DV}(\mathrm{n})$, and when all the questions have been asked we can check these against the decimal value obtained by the binary conversion of the 'yes/no' answers, SU (see Flowchart 5.8). The simplest thing for you to do now is to delete everything after line 260 and start entering from scratch again!

```
27G FRINT: "SCORE" : SU
OQ FFINT
```

SHO FOR $\mathrm{H}=1 \mathrm{TO} 5$

H）GOTO 40 D
3． HE CT N
330 FPINT，＂IJPEET HOT FOHAC．＂
4 EG FFIHT
410 PIJN
5 Fig IF－AIJT IN中
515 FN二1 IF IN中 $=$＂N＂THEN A $14=0$

53a Ev＝EW＋E？
545 PETUFN

1た可1日 $\mathrm{E}=1$
11 DAD［FTA EICYCE 2 A1
11015 ［JTA CAP． 172
$1192 E$ DATH TFARIH，EB
11 93g DATA PLAFAE， 175
1164 COTTA HOPSE， 12 ？
13 Din FOF $\mathrm{H}=1 \mathrm{TO} \mathrm{E}$

12G20 NEST H
13以


Flowchart 5．8 Matching the Decimal Value

This approach obviously saves a let of memory and time，as each array element takes up several bytes and must be located before it can be compared，so it is particularly useful where you are dealing with large
amounts of information. On the other hand, it does mean that you have to calculate the decimal equivalents of all of the bit patterns before you can use them, and it also gives you no clues when a complete match is not found. (Note that you cannotsimplytake the nearest decimal valuehere, as the decimal equivalent value of each correct answer depends on its position.) Of course you could do the calculations the hard way, but if you enter the bit pattern as a string, $1 \$$, then it is quiteeasy to convert it to the equivalent decimal value DV by comparing each single character slice MlD\$(I\$,N,l) with ' 1 ' and then adding on the value of the appropriate binary digit BD if a match is found.

```
2GHAC9 BO=1 : INFUT I$
25日15 FOF N=1 TOE
```



```
=[%/C[,
201936 B0=80+B0
2ด戸40}\mathrm{ NE:T H
20959 PRINT [U
2006[J PIN
```


## CHAPTER 6 <br> Making Your Expert System Learn for Itself

Although the＇expert＇systems described so far will function all right，they all require you to give them the correct rules on which to base their decisions in advance，which can be very tedious．

However，it is possible to construct an expert program which can learn from its mistakes and work out the decision rules for itself，provided that you can tell it when（although not where）it goes wrong．This is obviously an advantage if you are not altogether sure of the correct rules yourself anyway！In this case we start out with a series of features which should enable us to distinguish between the different objects，but without any pre－ defined yes／no pattern of these features（＇decision rule＇）to guide us．Instead we use the program itself to calculate what the pattern should be．

We will work with our familiar transport example and begin by setting up some variables．FE is the number of features to be considered（8）， $F E \$(n)$ is an array containing the names of these features，$F V(n)$ is an array which will hold the values which you give to each feature as input at any particular point（ 0 or 1 ），and $\mathrm{RU}(\mathrm{n})$ is an array which will hold the current overall values of the decision rule on each feature．

```
19 GOBLE 164%GE
10゙リリリ FE=8
```




```
15034 PEFC FEd[N;
LGG4E TEE%T N
11GDM [HATH WHEELS, HIHNS, EHNIFE,T
YRES,FHIL.S, WIN[DDWG, LHAIH,STEERIH
L
120゙5 PETINFH
```

Each feature is considered in turn（see Flowchart 6．1）．First the current feature value $\mathrm{FV}(\mathrm{n})$ for this cycle is set to zero，and then a＇yes／no＇input IN\＄is requested from the user on each point．If IN\＄is＇$Y$＇the feature value element $\mathrm{FV}(\mathrm{N})$ is set to I ；otherwise it remains set at zero．This will produce a pattern which describes the object as＇ 0 ＇and＇ 1 ＇in array FV（n）．


Flowchart 6．1 Learning to Distinguish Between Two Objects

```
6可 FIOP N=1 TO FE
76 FVKN \=ら
G0 PPINT FE$CN:;" ";
```



```
10H PRINTT IN$,
115 IF IN{界="''" THEN FVKN\rangle=1
12@ NE:T N
```

Before you start a decision variable DE is set to zero．This is recal－ culated as the sum of the current value of DE ，plus each of the feature values $\mathrm{FV}(\mathrm{N})$ entered，multiplied by the current decision rule values RU（N）．

```
125 [DE=ら
130 FDF N=1 TO FE
15f] [PE=[DE+F\(N)FFU(N)
165 NE%T N
17可 PRINT "DE= ":DE
```


## Which is which？

To start with we will consider the simplest situation where there are only two possibilities－a bicycle or a car．Initially we make the distinction between these quite arbitrarily by saying that if the final value of DE is equal to or greater than 0 then it is a bicycle，whereas if DE is less than 0 then it is a car．It does not really matter that this is not actually true as the system will soon correct itself．When the program has made a decision on the basis of the value of DE it requests confirmation（or otherwise）of the result．

```
180 IF CE>=@ THEN PRINT"IS IT A
BICYCLE ";:INPUT IN%:GOTO 200
190 IF [DE<O THEN PRINT"IS IT A C
HR ";:INPUT IN*:GIJTO 22פ
```

Three possible courses of action may be taken according to whether or not the computer＇s decision was correct．If it was correct then effectively no action is taken（a weighting variable WT is set to zero），and the program loops back for another try．If DE was $>=0$ but the computer was wrong， then the weighting variable WT is set to minus one，whilst if DE was $<0$ but the computer was wrong then WT is set to plus one．

## 206 IF IN海＂＇Y＂THEN WT＊ด：GOTO 24 <br> ด <br> 210 WT＝－1：GOTO 240 <br> Z？ด IF IN\＆ェ＂＂く＂THEN WT＝ด：IOTO 24 <br> ज

230 WT $=1$
The effect of the weighting variable is to modif $y$ the values in the rule array $\mathrm{RU}(\mathrm{N})$ ，pulling them down when they are too high，and pulling them up when they are too low．

240 FOR H＝1 TO FE
350 RU（N）＝FU（N）＋FVCN）：WT
260 PEINT RU（N），
2？ CH WT H
206 PRINT
2ロब חIOTM 60

The way the system operates is best seen by a demonstration．Type RUN and then follow this sequence of entries．（Note that the punctuation has been designed to give a screen format which clearly indicates the relationship between your input values and the decision rule values．）

First of all enter these values：

| WHEELS $Y$ | WINGS $N$ | ENGINE N | TYRES Y |
| :--- | :--- | :--- | :--- |
| RAILS $N$ | WINDOWS $N$ | CHAIN Y | STEERING Y |

The program will return with a decision value DE of zero, as this is the initial value and no modifications have yet taken place:
$\mathrm{DE}=0$

As DE is 0 , the system assumes that this is a bicycle and asks for confirmation, to which the answer is of course 'yes'.

## IS IT A BICYCLE? Y

The print-out of the contents of the rule array $\mathrm{RU}(\mathrm{n})$ shows that these have not changed fromzero as the correct answer was, by pure chance, obtained:
0
0
0
0
0
0
0
0

Now try entering this sequence, which describes a car:

| WHEELS $Y$ | WINGS $N$ | ENGINE Y | TYRES Y |
| :--- | :--- | :--- | :--- |
| RAILS $N$ | WINDOWS Y | CHAIN $N$ | STEERING Y |

DE is still zero, so the wrong conclusion is reached and the wrong question is asked, to which the answer must be 'no':

```
DE=0
IS IT A BICYCLE?N
```

Now, as a mistake was made, the decision rule is modified by subtracting one from each value in the rule array where a 'yes' answer was given. The contents of the rule array thus become:

| -1 | 0 | -1 | -1 |
| :--- | :--- | :--- | :--- |
| 0 | -1 | 0 | $-I$ |

If you once more enter the values which describe a car, the program will come up with the correct answer:

| WHEELS $Y$ | WINGS $N$ | ENGINE Y | TYRES Y |
| :--- | :--- | :--- | :--- |
| RAILS N | WINDOWS Y | CHAIN N | STEERING Y |

```
DE=-5
IS IT A CAR?Y
```

| -1 | 0 | -1 | -1 |
| :--- | :--- | :--- | :--- |
| 0 | -1 | 0 | -1 |

Bef ore you feel too pleased with yourself, try giving it the values for a bicycle again, which it will get wrong!

| WHEELS Y | WINGS N | ENGINE N | TYRES Y |
| :--- | :--- | :--- | :--- |
| RAILS N | WINDOWS N | CHAIN Y | STEERING Y |
| DE=-3 |  |  |  |
| IS IT A CAR?N |  |  |  |


| 0 | 0 | -1 | 0 |
| :--- | :--- | :--- | :--- |
| 0 | -1 | 1 | 0 |

However the positive features which are common to the bicycle and the car are now automatically increased by one, so that if you repeat this last sequence it will now produce the correct conclusion:

| WHEELS Y | WINGS N | ENGINE N | TYRES Y |
| :--- | :--- | :--- | :--- |
| RAILS N | WINDOWS N | CHAIN Y | STEERING Y |

$D E=1$
IS IT A BICYCLE? Y

| 0 | 0 | -1 | 0 |
| :--- | :--- | :--- | :--- |
| 0 | -1 | 1 | 0 |

The situation has now stabilised and the program will always recognise both car and bicycle correctly every time you enter the features which describe them:

| WHEELS Y | WINGS N | ENGINE Y | TYRES Y |
| :--- | :--- | :--- | :--- |
| RAILS $\mathbf{N}$ | WINDOWS Y | CHAIN $\mathbf{N}$ | STEERING Y |

DE $=-2$
IS IT A CAR? Y

| 0 | 0 | -1 | 0 |
| :--- | :--- | :--- | :--- |
| 0 | -1 | 1 | 0 |

Notice that the final value of DE for a bicycle is I, and for a car-2. If you look at the rule array values, you will see that these correspond in both number and position to the unique features which distinguisht hese objects (CHAIN for bicycle, and ENGINE and WINDOWS for car).

## A wider spectrum

Although you have now managed to teach your computer somet hing, it is not exactly earth-shattering for it to be able to distinguish bet ween only two objects. Let's expand the system to deal with a wider spectrum of possibilities (see Flowchart 6.2). To start with we need to define the


Flowchart 6.2 Learning the Rules for a Wider Spectrum of Possibilities
number of objects we wish to be able to recognise OB, name them as DATA which we READ into a new array OB\$(OB), change our decision rule array into a two-dimensional form, $\mathrm{RU}(\mathrm{FE}, \mathrm{OB})$, which can hold rules for each of the objects separately, and set up a decision array DE(n) to hold decision values for each object.

```
10 GOSIJB 10EFB
10009 FE=8:0E=5
10@10 DIM FE&(FE),F\(FE),PINFE,D
E`,UE$(UE),DE(UE)
1GD20} FOR N=1 TO FE
1053O READ FE$N
10040 NE:KT N
1005G FOR N=1 TO OB
1GEEG RERD UE$(N)
10070 NEXT H
1100g CATF WHEELE,WINGGS,ENGIHE,T
YPES, RAIL.S, WIMCOWS, CHAIH, STEERING
11以10 [ATA BICYCLE,CAF, TRAIN,PLR
HE. HORSE
1%000 FETI.JPN
```

Rather than just having a single decision variable DE ，we need here to determine a decision value for each object each time．In each cycle we must first set $D E$ to zero，and then zero every element in the decision array $\mathrm{DE}(\mathrm{n})$ so that we start with a clean slate for every object．

```
20 DE=0
35 FOR N=1 TO OB
45 DE(N)=E
55 NEXT N
```

The values for each feature are then entered in exactly the same way as before．

```
5 0 ] ~ F O R ~ N = 1 ~ T O ~ F E
70 FW(N)=0
90 PRINT FE#(N`;" ";
90 IN% =1NKEY*:IF IF性="" THEN 90
10W PRINT IN&,
110 IF IN泫="Y" THEN FV/N`=1
12O NE%T N
```

Each element of the decision array $\mathrm{DE}(\mathrm{n})$ is now updated according to the status of the entered values $\mathrm{FV}(\mathrm{n})$ and the contents of the appropriate rule array element $\mathrm{RU}(\mathrm{n}, \mathrm{m})$ ．

130 FOF $\mathrm{H}=1 \mathrm{~T}$ TI FE
145 FOP $14=1$ TO OB

1E日 HENTM，M

We now need to look to see if any of the decision values for any of the objects $\mathrm{DE}(\mathrm{n})$ are greater than or equal to the overall decision value DE ． If this is true，we set a＇top score＇TS variable equal to the number of the object producing the best match．

```
IPG FDR N=1 TIJ DE
180 IF [JE(M)>=[DE THEN OE=0E(N):TS=H.
199 NE\T &
```

The best guess of the system is that this is the correct answer，so once again it asks for confirmation，and simply returns for a new input without making any changes if the answer was correct．

```
200 FRINT "WFS IT ":OE*CTS>:" ":
```

210 IH
25 FRINT IN
ここめ IF IN\&="个" THEH 2ら

If this was not the correct answer，the names and numbers of all the objects are printed out and you are asked for the number of the correct answer CR．（The limitations on CR prevent you crashing the program by entering an illegal value．）

230 FOR $\mathrm{H}=1$ TO DE
24日 FRINT N，OE\＆CN？
ZSB NEMT H
260 FRINT＂WHICH WRS 1T＂：
270 IHFUT ER：IF CREI OF CP\％THE
H 270

A check is now made to see if the decision value for each object $\mathrm{DE}(\mathrm{n})$ is greater than or equal to the overall decision value DE and whether the object being considered is not the correct answer．If both of these are true then the rules are updated again by subtracting the correct feature values $\mathrm{FV}(\mathrm{n})$ to bias in favour of the correct answer．

```
20G FOR N=1 TO OB
290] FRIMT [DE(N);DE;CR
300 IF [DE(N`)=[JE FNHD NK`CO THEN
FOR M=1 TO FE:PIOM,N)=FI(M,N)-FW
(M):HE%T M
310 NE%T H
```

Now the correct feature values $\mathrm{FV}(\mathrm{n})$ are added to the rule array for the correct object，to bias in the opposite direction．

```
325 FOF M=1 TOI FE
33D F|J(M,CR)=F|KM,IP +F\/6M%
3 4 0 ~ स E : K T ~ M ~
```

Finally the statuses of the rule arrays are printed out so that you can see what is happening.

```
35 FOF M=1 TO OR
3Eด FOJF N=1 TO FE
30}\mathrm{ PFINT F|{N,M);
3日G NEMT N
390 PRINT
455 HEXT M
410 GOTO ?0
```

Once again a demonstration is the best way to understand what is happening so enter the following sequence:

| WHEELS Y | WINGS N | ENGINE $N$ | TYRES Y |
| :--- | :--- | :--- | :--- |
| RAILS N | WINDOWS N | CHAIN Y | STEERING Y |

The program will come back with the erroneous conclusion that it was a horse, so you must tell it that this was wrong, when it will ask you for the correct answer (bicycle $=1$ ):

## WAS IT HORSE N

1 BICYCLE
1 CAR
3 TRAIN
4 PLANE
5 HORSE

## WHICH WAS IT I

The statuses of the various decision and rule arrays are now printed out for your information (note that the labels shown here are not included on the screen).

| $(\mathrm{DE}(\mathrm{N}))$ | $(\mathrm{DE})$ | $(\mathrm{CR})$ |
| :--- | :--- | :--- |
| 0 | 0 | 1 |
| 0 | 0 | 1 |
| 0 | 0 | 1 |
| 0 | 0 | 1 |
| 0 | 0 | 1 |


| 1 | $\mathbf{0}$ | 0 | $\mathbf{1}$ | 0 | 0 | $\mathbf{1}$ | $\mathbf{1}$ | (bicycle) |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| -1 | 0 | 0 | -1 | 0 | 0 | -1 | -1 | (car) |
| -1 | 0 | 0 | -1 | 0 | 0 | -1 | -1 | (train) |
| -1 | 0 | 0 | -1 | 0 | 0 | -1 | -1 | (plane) |
| -1 | 0 | 0 | -1 | 0 | 0 | -1 | -1 | (horse) |

A B C
D E F
G $\quad \mathrm{H}$

| $(\mathrm{A}=$ wheels | $\mathrm{B}=$ wings | $\mathrm{C}=$ engine | $\mathrm{D}=$ tyres |
| :---: | :--- | :--- | :--- |
| $\mathrm{E}=$ rails | $\mathrm{F}=$ windows | $\mathrm{G}=$ Chain | $\mathrm{H}=$ Steering) |

If you look closely you will see that the features which have caused alterations in the rule arrays are wheels, tyres, chain and steering - all features which we defined as part of a bicycle and not found in a horse. In addition, you will see that the values for these features in the bicycle rule array are now all plus one, whilst the values for these features for all the other objects are now all minus one.

Now give it the features of a car, which it will think a bicycle, and then correct it. Notice that the rule arrays for bicycle and car are now amended to take into account the new information.

| WHEELS Y | WINGS N | ENGINE Y | TYRES Y |
| :--- | :--- | :--- | :--- |
| RAILS $N$ | WINDOWS Y | CHAIN N | STEERING Y |

## WAS IT BICYCLE N

| $\mathbf{1}$ | BICYCLE |
| :--- | :--- |
| $\mathbf{2}$ | CAR |
| $\mathbf{3}$ | TRAIN |
| $\mathbf{4}$ | PLANE |
| $\mathbf{5}$ | HORSE |

WHICH WAS IT 2

| 3 | 3 | 2 |
| ---: | ---: | ---: |
| -3 | 3 | 2 |
| -3 | 3 | 2 |
| -3 | 3 | 2 |
| -3 | 3 | 2 |


| 0 | 0 | -1 | 0 | 0 | -1 | 1 | 0 | (bicycle) |
| ---: | ---: | ---: | ---: | :--- | ---: | ---: | ---: | :--- |
| 0 | 0 | 1 | 0 | 0 | 1 | -1 | 0 | (car) |
| -1 | 0 | 0 | -1 | 0 | 0 | -1 | -1 | (train) |
| -1 | 0 | 0 | -1 | 0 | 0 | -1 | -1 | (plane) |
| -1 | 0 | 0 | -1 | 0 | 0 | -1 | -1 | (horse) |


| A | B | C | D | E | F | G | H |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |


| $(\mathrm{A}=$ wheels | $\mathrm{B}=$ wings | $\mathrm{C}=$ engine |
| :--- | :--- | :--- |
| $\mathrm{E}=$ rails | $\mathrm{F}=$ windows | $\mathrm{G}=$ chain |

Next give it a plane, which it decides is a car, and correct it again.

| WHEELS Y | WINGS Y | ENGINE Y | TYRES Y |
| :--- | :--- | :--- | :--- |
| RAILS N | WINDOWS Y | CHAIN N | STEERING Y |

## WAS IT CAR N

1 BICYCLE
2 CAR
3 TRAIN
4 PLANE
5 HORSE

## WHICH WAS IT 4

And now a train, which it still gets wrong!

| WHEELS $Y$ | WINGS N | ENGINE Y | TYRES N |
| :--- | :--- | :--- | :--- |
| RAILS $N$ | WINDOWS Y | CHAIN N | STEERING $N$ |

## WAS IT PLANE N

1 BICYCLE
2 CAR
3 TRAIN
4 PLANE
5 HORSE

## WHICH WAS IT 3

And finally a horse, which comes out as a plane!

| WHEELS N | WINGS N | ENGINE N | TYRES N |
| :--- | :--- | :--- | :--- |
| RAILS N | WINDOWS N | CHAIN N | STEERING Y |

WAS IT PLANE N

1 BICYCLE
2 CAR

```
3 TRAIN
4 PLANE
5 HORSE
```


## WHICH WAS IT 5

If you continue to feed your expert information, eventually it will get the right answer every time. How long this will take depends upon the extent of the differences between the features of the objects, and on the order in which the objects are presented to the expert. Be warned that it can take a long time before it becomes infallible. Here is one sequence which eventually came out right every time.

| plane (train) | car (plane) | bicycle (YES) |
| :--- | :--- | :--- |
| car (YES) | plane (car) | plane (YES) |
| horse (YES) | plane (bicycle) | car (plane) |
| plane (car) | plane (car) | car (plane) |
| car (YES) | plane (car) | plane (YES) |
| car (YES) | plane (YES) | horse (YES) |
| bicycle (YES) | train (car) | train (YES) |
| bicycle (YES) | car (plane) | car (YES) |
| plane (car) | plane (YES) | car (plane) |
| car (YES) | plane (YES) | car (YES) |
| bicycle (car) | car (YES) | plane (YES) |
| train (YES) | horse (YES) | bicycle (YES) |

To see the final state of the rule array when the ultimate state is reached, you can stop the program and then type GOTO 350 as a direct command. As the final scale of values ranges from +6 to -2 , you should not be surprised that it took a long time to get there.

| 1 | 0 | -1 | 1 | 0 | -2 | 3 | 0 | (bicycle) |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | :--- |
| -1 | 4 | 1 | 0 | -1 | 1 | -2 | 0 | (car) |
| 0 | -1 | 1 | -2 | 2 | 1 | -1 | -2 | (train) |
| -2 | 6 | 0 | 0 | -1 | 0 | -2 | -2 | (plane) |
| -1 | 0 | 0 | -1 | 0 | 0 | -1 | 0 | (horse) |
| A | B | C | D | E | F | G | H |  |
| (A=wheels | $\mathrm{B}=$ =wings | $\mathrm{C}=$ engine | $\mathrm{D}=$ tyres |  |  |  |  |  |
| $\mathrm{E}=$ rails | F =windows | $\mathrm{G}=$ chain | $\mathrm{H}=$ steering) |  |  |  |  |  |

Of course, in a real application of such an expert system you could feed it a
mass of collected information and conclusions on a subject area and then leave it alone to digest this and to come up with the rules in its own good time. As these rules are stored in arrays you could easily write a routine to save these for re-use later.

## CHAPTER 7

## Fuzzy Matching

Computers are totally logical but our own memory banks are much more unreliable, which can lead to problems when you are trying to recover information on a particular subject. Forexample, English is a very variable language and there are frequently alternative spellings of the same (or very similar) surnames, which can cause difficulties.

One way around this problem is to try to match the sound of the word, rather than the actual letters in it, by means of 'Soundex Coding', which was originally de veloped to ass ist processing of the 1890 census in the USA. This method of coding ensures that similar-sounding words have almost the same code sequence.

The rules for coding a word are as follows:
I) Always retain the first letter of the word as the first character of the code.

From the second letter onward:
2) Ignore vowels (a, e, i, o, u).
3) Ignore the letters w, y, q and h.
4) Ignore punctuation marks.
5) Code the other letters with the values $1-6$ as follows:

| Letters | Code |
| :--- | :--- |
| bfpv | 1 |
| $\mathrm{cgjks} \times z$ | 2 |
| dt | 3 |
| l | 4 |
| mn | 5 |
| r | 6 |

6) Where adjacent letters have the same code only the first one is retainec
7) If length of code is greater than four characters then take first four only.
8) If length of code is less than four characters then pad out to four characters with zeros.

To make this clear here are some examples of Soundex Coded names:

BRAIN - B650
( B is retained, R is $6, \mathrm{~A}$ and I are dropped, N is 5 and a zero is added to pad out the code.)

CUNNINGHAM - C552
( C is retained, U is dropped, both N s are represented by the single code $5, \mathrm{I}$ is dropped, the third N is represented by $5, \mathrm{G}$ is $2, \mathrm{H}$ and A are dropped, and M is coded as 5 - but the resulting code (C5525) is truncated to four characters.)

GORE - G600
( G is retained, O is dropped, R is 6 , E is dropped and zeros are added to pad the code.)

IRELAND - 1645
(I is retained, R is $6, \mathrm{E}$ is dropped, L is $4, \mathrm{~A}$ is dropped, N is 5 and D is $3-$ but the resulting code (16453) is truncated to four characters.)

SCOT - S230
( S is retained, C is 2 , O is dropped, T is 3 and zero is added to pad the code.)
If your name is full of vowels and other rejected letters, then you will find that your code is somewhat abbreviated!

HEYHOE - H000
( H is retained, all the other letters are rejected (!), and the code is filled up with zeros.)

## Coding routine

To save all that brainwork, let's develop a program which allows you to input a word in English and output it in Soundex Code (see Flowchart 7.1). The first thing to do is to jump to a set-up routine which reads each group of the retained letters into one element of a Soundex Code string array SC\$(n). (Note that the letters are arranged so that they are in the array element corresponding to their code value.)


Flowchart 7．1 Producing a Soundex Code

```
10GOS|P 10ほうに
10050 [IM SCW(E)
110MD ORTA RFFY,CK.JKS%'Z,DT,L,MN,R
120日0 FOR N=1 TO G
12g15 PEAD SC*N)
129%O REXNTN
13000 RETIJPH
```

We can now input the word to be converted，INS，and，to begin with， make the coded version of this，COS，the first letter of that word （following the first rule above）．

## 10ロ INPUT IN\$ <br> 

We now need to check the other letters of the word, 2 TOLEN(IN\$), in turn after first making a temporary string TM\$ equal to the current letter to be translated.

## 120 FOP $\mathrm{N}=2$ TO L.ENGIN期? <br> 

As conversion to the code numbers will be required at various points in the final problem, we will set up this process as a subroutine at line 1000.

## 140 EOSIJ 1000

We have to check TM\$ against each individual letter in each group of letters $\operatorname{SC}(n)$ to find a match. To check each letter group, we have to go round six times, making a search string SE\$ the current Soundex Code group, and using INSTR to check each letter in the group against TM\$ inturn.

```
100G FOR P=1 TU E
1010 SE$=SC&(P)
1020 SP=INSTP(1,SE&,TM%)
```

When the INSTR check has been made, we have to determine whether a match has been found to any of the Soundex groups, and if so, to which group. If no match is found then SP will be set to zero. If a match is found then SP will be set to $M$ which will point to the value of the code group matched.

If a match is found, $\mathrm{SP}>0$, then we convert the numeric value of the loop scanning the code groups P to a string TM\$ which replaces our original temporary string. (The STR\$ command converts a number into a string, but we also need to use RIGHT\$ as STR\$ automatically adds a space on to the front of the number string.)



If no match is found in that group, we have to check the next group.

## 1040 NEST F

If no match is found at all, then TM\$ must contain one of the characters to
be ignored．So we reset TM $\$$ empty［ $\$=\times "]$ and RETURN．

## 1050 TM化＝＂＂ <br> 1 160 FETUJRN

We can now make the coded string CO\＄equal to the original coded string plus the newly converted character TM\＄．

## 

Now we loop back to deal with the next character in IN\＄．

## 18日 NENT H

When the end of IN\＄is reached，we print out the input IN\＄and the entire coded string CO\＄before going back to 100 for another input．

```
210 PFINT PRINT "NANE", "OCDE":FR
INT IN&.CO$
320 F..OTU 100
```

If you input the name STE VEN you will nowget the code S3I5 which is correct．However，if you try BRAIN or CUNNINGH AM you will get the codes B65 and C55525 respectively．The code for BRAIN is too short and needs padding out with zeros，and the code for CUNNINGHAM is too long and the same codes are repeated one after another for the letter $\mathbf{N}$ ．

## Dealing with the details

To solve the problem of the repetition of the same code for adjacent letters， we need to keep a record of the last temporary string LT\＄．We need to make LT\＄the code of the first character in IN to start with，so that the initial letter is not repeated．As we go through the FOR－NEXT loop．we must compare LT\＄with TM\＄，and if they are the same we must not add TM \＄to CO\＄．Otherwise we need to make LT\＄the latest TM\＄．



```
150 IF TH性=L.T$ THENGITG 1EG
160 I.T$=T川各
```

Now we can sort out the problem oft he codebeing too short．First of all we check the length of the string $\operatorname{LEN}(\mathrm{CO} \$)<4$ ．If it is too short，we add three zeros on to the end and then use LEFT\＄to cut the string back down to the correct size（four characters）．


Flowchart 7.2 Dealing with the Details



Finally, if the string is too long then we cut it down to size with LEFT\$(CO\$,4) again.

##  CCDF, 4 )

## Matchmaking

Now that we have a reliable method of producing the Soundex Codes，let＇s give it something to work on．The first task is to read a list of names out of DATA statements into a name string array NA\＄（n）． Our demonstration list only consists of eighteen names－if you want more，a quick flick through your local telephone directory should soon solve that problem！Note that the number of words is also stored as NW．

```
10015 トMN=17:DIM NF$(NW)
1015 [ARTR RERRHAI, RERRHFIN, RERA
MS, R[ORM, R[SAMS, A[DORMS, A[JRMSON, RLA
14. RLLFN, PLI.EN
1102g ['ATA RNTHRN'Y', RNTHONY', RNTDDN
'r.. AITTPORUS, APPEPLEY., APPLEEEE, FPP
LEE'r'.RPPLEFORD
12630 FOF N=0 TO 1?
12940 PERD NR&(N)
1295g NE:'T N
```

The whole idea of matching with Soundex Codes relies on the fact that you use the Soundex Code to make the match before printing the possible words．We therefore have to find the codes for each of the names from the DATA and put these codes into an equivalent string array $\mathrm{NC} \$(n)$ ．The routine to find the Soundex Code is virtually identical to the one used to find the code of an input，as described above．

```
1@g%a [IM NC&(NW)
1:\DE@ PRINT:PRINT "KRME","CDDE":
PRIFIT
12070 FOR 0=0 TO NW
129:09 PFIF,T NR&C(0),
```



```
*-GOSLJB 1000:LT每与TM年
1210日 FOR N=2 TO LEN(NR&(O))
12110 TM利\0$(NA$(0),N,1)
12120 C.0SUS 1000
12139 IF TM%=LTक THEN HE\T N:GOT
0 121?0
12140 LT出利井
12150 CDक=С味+TM复
12160 NEYT N
1212. IF LEN(CO*)<4 THEN CD$=CO$
+"0日ด":CO$=LEFT%(CD$,4)
```

```
1219历 IF LEN( CO'% ):4 THEN CO$=LEF
T$(C05,4)
12190 PRINT C[j% 
12200 NC%( ) ==CO%
12210 NEXT Q
```

If you RUN this now，you will see all the codes for the DATA produced before the input request．

| NAME | CODE |
| :--- | :--- |
| ABRAHAM | Al65 |
| ABRAHAMS | Al65 |
| ABRAMS | Al65 |
| ADAM | A350 |
| ADAMS | A352 |
| ADDAMS | A352 |
| ADAMSON | A352 |
| ALAN | A450 |
| ALLAN | A450 |
| ALLEN | A450 |
| ANTHANY | A535 |
| ANTHONY | A535 |
| ANTONY | A535 |
| ANTROBUS | A536 |
| APPERLEY | A164 |
| APPLEBEE | A141 |
| APPLEBY | A141 |
| APPLEFORD | A141 |

The only thing we need to do now is to find which codes of these names match the code of your input and then to print out these names with a FOR－NEXT loop．

```
245 PRINT
250 FOR N=麦 TO NW以
26[5 IF COw-NC$(H) THEN FPINT NA&
〔N%NC(W@N)
27ら NEXT N
```

This will only print words with exactly matching Soundex Codes．For
example, if you try entering the name APPLEBE you will get the following response:
? APPLEBE

| NAME | CODE |
| :--- | :--- |
| APPLEBE | Al41 |
|  |  |
| APPLEBEE | Al4! |
| APPLEBY | Al4! |
| APPLEFORD | Al4! |

Although APPLEBE (one E at the end!) is not present in the DATA, we have found APPLEBEE and APPLEBY, as well as APPLEF ORD (where the interesting sound at the end has been chopped off).


Flowchart 7.3 Partial Matching

## Partial matching

Notice that on the other hand APPERLEY has been rejected, even though it sounds quite similar at first. It would therefore be useful if we could also print out partial matches.

This can easily be done by adding an extra FOR-NEXT loop, which compares a decreasing section of the input LEFT\$(CO\$,M) with decreasing lengths of the stored codes LEFT\$( C\$(N),M) (see Flowchart 7.3).

```
230 FOR M=4 TO 1 STEP -1
240 PRINT.PRINT M;"CHRFACTERS IAR
TCH":PRINT
26日 IF LEFT$(CO&,M)=L.EFT&(NC$(N)
M> THEN PRINT NRW(N),NCG(N)
290 PRINT PRINT "PRESS KEY TO CO
MT {NUE"
290 IN4=INKE'Y年:IF IN&="" THEN 29
\square
300 PRINT PRIMT
31D NEKT M
```

If you now try APPLEBE you can see the whole range of possibilities.

## ? APPLEBE

| NAME | CODE |
| :--- | :--- |
| APPLEBE | Al41 |

4 CHARACTERS MATCH
APPLEBEE ..... Al4l
APPLEBY ..... Al4l
APPLEFORD ..... Al41
PRESS KEY TO CONTINUE
3 CHARACTERS MATCH
APPLEBEE ..... Al41
APPLEBY ..... A141
APPLEFORD ..... A141
PRESS KEY TO CONTINUE
2 CHARACTERS MATCH
ABRAHAM ..... A165
ABRAHAMS ..... A 165
ABRAMS ..... A 165
APPERLEY ..... Al 64

| APPLEBEE | Al41 |
| :--- | :---: |
| APPLEBY | Al41 |
| APPLEFORD | Al41 |
| PRESS KEY TO CONTINUE |  |
|  |  |
| I CHARACTERS | MATCH |
| ABRAHAM | A165 |
| ABRAHAMS | Al65 |
| ABRAMS | A165 |
| ADAM | A350 |
| ADAMS | A352 |
| ADDAMS | A352 |
| ADAMSON | A352 |
| ALAN | A450 |
| ALLAN | A450 |
| ALLEN | A450 |
| ANTHANY | A535 |
| ANTHONY | A535 |
| ANTONY | A535 |
| ANTROBUS | A536 |
| APPERLEY | A164 |
| APPLEBEE | A141 |
| APPLEBY | A141 |
| APPLEFORD | A141 |
| PRESS KEY TO CONTINUE |  |

## CHAPTER 8

## Recognising Shapes

We normally recognise objects using our senses of sight, sound, taste and feel, whereas of course our basic computer can only obtain information through the keyboard. Whilst it is possible to produce sensors which can be interfaced with your machine to give it another view of the outside world, constructing these requires a reasonable amount of electronic and mechanical knowledge and skill. We will makedo instead with a simulation of the action of a light sensor to illustrate how shapes can be recognised.

Let us think for a start about three simple shapes - a vertical line, a square, and a right-angled triangle.
We can recognise these shapes by looking at the pattern they make on an imaginary grid and deciding whether or not there is a point set at each X and $Y$ coordinate.
In the case of a line only the first X coordinate is used, but all of the Y coordinates. A square is a little more complicated, as all the X coordinates on Y rows I and 8 are set, and from Y rows 2 to 7 only the first and last X points are set. Finally, a triangle is even more complicated, as the slope is produced by incrementing the X axis each time

Table 8.1 Decimal Values of Shapes Described in Binary Form

| Y row | line | square | triangle |
| :--- | :--- | :--- | :---: |
| 1 | 1 | 255 | 1 |
| 2 | 1 | 129 | 3 |
| 3 | 1 | 129 | 5 |
| 4 | 1 | 129 | 9 |
| 5 | 1 | 129 | 17 |
| 6 | 1 | 129 | 33 |
| 7 | 1 | 129 | 65 |
| 8 | 1 | 255 | 255 |

One obvious way to describe these particular figures would be to represent each point by a single bit and to produce a decimal value for each row, in the same way as we did before when we were looking at expert systems (see Table 8.1). In fact this type of approach is used to produce the characters which you see on your screendisplay, the formats for which are
stored in memory in just this form. For example Figure 8.1 shows how the letter ' $A$ ' is built up.

There are now machines availa ble(Optical Character Readers) which can reverse this process. They actually 'read' a printed page by scanning the paper in a grid pattern and measuring whether or not light is reflected at particular coordinates.


Figure 8.1 Forming the Letter ' $\mathbf{A}$ '

What they actually take in will be a pattern of 'yes' and 'no' for each coordinate, and of course this must then be decoded and compared with the patterns for known shapes. The most obvious way to make this comparison would be to consider every point in turn as a binary digit and then to convert each row back to a decimal value which could be compared with a table of known values. However this has the disadvantage that we must actually check every individual point on the grid ( 64 points).

## A branching short cut

A quicker approach relies on the fact that each character can actually be detected by looking at only a much smaller number of critical features of the pattern. For example, Figure 8.2 gives a decision tree which will find all


Figure 8.2a Decision Tree for Alphabet


Figure 8.2b
the capital letters of the alphabet using only 12 points (see Figure 8.3), and it is not even necessary to check all 12 in any particular case. If you follow each of the routes, you will see that the maximum number of steps to be


Figure 8.3 Points Used in Decision Tree
followed is seven, and that most letters are found in less than five steps (Table 8.2). This must obviously be quicker than comparing all 64 points!

Table 8.2 Numbers of Steps Required for Recognition of Each Character

```
3 steps - I, D
4 \text { steps - L, J, C, G, O, W}
5steps - S, A, Q, R, T, F, U, space
6 \text { steps - P, V, Y, H}
7 steps - B, M, N, E, K, X, Z
```

To demonstrate how this approach works，we will simulate the action of the scanning head by producing a grid on the screen，on which you can construct characters．

The text screen start address 1024 is defined as a variable TS as it is used frequently．The screen is cleared and a red area $6 \times 8$ blocks is set up in the top lef thand corner by POK Eing CHR\＄（191）．A yellow（CHR\＄（159）） $5 \times 7$ grid is then superimposed on this to mark the actual working area（of course there must be a margin around the edge so that characters do not merge）．

```
10 GUSUE 10000
100@G TS=1924
1200D CLS
12D10 FOF N=1 TO 10
12ดこD PRIHT
1203D HES价 N
13100D FOR ン叮TO 曰
131310 FOF )'=i' TO &
13020 POKE TS+Y名隹32 ),191
1303D I'ENT%
13040 NEUT Y
13D50 FOP )}=1\mathrm{ TO 5
13DGD FOP Y=1 TO ?
1 33PD POKE TS+Y+{'r'32 2,159
13DED IHENT 分多
13D90 : =1: Y=1
1310D PETUFN
```

flashing cursor is now produced to show your position． CP is the current position on the text screen，TS，the current colour of which is saved as CC by PEEK ing this position．A different coloured block CC +32 is POKEd into place and thenthe original colour CC POK Ed back，so that there is no lasting effect．If no key is pressed，the program loops back and checks again：when a key is pressed，the ASCII value of $A \$$ is taken as A．

2月 A 事＝INKE

$\mathrm{P}, \mathrm{EC}+\mathrm{ZC}: \mathrm{POKE} \mathrm{P}, \mathrm{CC}$

＊）

The X and Y coordinates are updated according to movement of the cursor keys（ $A=9,8,10$ or 94 ），and if the spacebar（ $A=32$ ）is pressed the colour of the current position is set to dark blue，CHR\＄（175）．If you make a mistake，
pressing＇ X ＇erases the current position by resetting the colour to yellow， CHRS（159），or the CLEAR key（ $\mathrm{A}=12$ ）jumps to the set up routine and erases all the current grid．Pressing RETURN（ $A=13$ ）leads to the decoding routine，or else the program loops back to the keycheck．

```
5 0 ~ I F ~ R = 9 ~ T H E N ~ X = \% + 1 ~
60 IF R=OS THEN K=X-1
70 IF A=1ด THEN 'O=Y+1
90 IF R=94 THEN Y=Y-1
9% IF AF=%2 THEN POKE TS+Y+(Y*32)
,175
1&行 IF R然="S" THEN POKE TS+%+CY*
3.2),159
110 IF R=12 THEN EOSLIE 13000
120 IF R=13 THEN 200口
130 IF K<1 THEN 
140 IF X>S THEH X=S
150 IF `队 1 THEH Y=1
160 IF Y>7 THEH Y=7
170 GOT\200
```

Limits must be set to prevent the cursor wandering off the $5 \times 7$ grid area．

```
13ด IF %1 THEN:=1
140 IF %S THEH:%=5
tGG IF %<< THEN ' }=
16@ IF '%'? THEH\ '=?
```

The decision tree is held in a series of linked arrays where $N B$ is the nu mber
 to be checked next， $\mathrm{C} 2(\mathrm{n})$ the Y coordinate to be checked next， $\mathrm{N}(\mathrm{n})$ the next element to use if theansweris＇no＇，and $Y(n)$ the nextelementtouseifthe answer is＇yes＇．

```
1109ロ 1-5:53
```



```
N+4E%,%(NE)
11月20 FDR r&=1 TO HE:
```



```
H2,Y(N)
11%40 HEYT H
```

each branch point), as this makes it easy to enter and to edit out any mistakes.

| 14915 | [UATA , 1, 1, 2, 19 |
| :---: | :---: |
| 1,400 | DATA .1, 5, 3, 19 |
| 14035 | [AT'A , 3, 2.4.3 |
| 14045 | URTA , 5, 1, 5, 8 |
| 14595 | DATA 3 |
| 14050 | DATA " |
| 14975 | [PATF " 8 " |
| 14090 | OATA ".J" |
| 14990 | [ATA "I |
| 14190 | DATA , 5, 4, 11,14 |
| 14110 | DATA,5, 3, 12,13 |
| 14129 | OATA "C" |
| 14130 | DATA "E" |
| 14140 | [ARTA , 5, $7,19,15$ |
| 14150 | DATA , 2, 4, 17,16 |
| 14169 | DATA "R". |
| 14175 | DATE "O." |
| 14180 | DATH "O" |
| 14190 | [ARTA , 5, 1, 20, 29 |
| 14290 | DATA , 5, 4, 21,29 |
| 14219 | DATA , 5.3.27,22 |
| 14290 | DATA , 5, 7, 23,26 |
| 14230 | DATA , 5,5,24,25 |
| 14240 | [ ${ }^{\text {STTA }}$ "P" |
| 14259 | DATA "E". |
| 14260 | DATA "R". |
| 14270 | DATA "L", |
| 142 '90 | DATA "L" |
| 14295 | [DATA , 5, 7, 45,32 |
| 143 ¢ 0 | DATA , 2, 6, 31,44 |
| 14310 | [,ATA , 5, 3, 3, 2,35 |
| 14320 | DATA , 1, 5, 33,36 |
| 14330 | DATA , 3, 1, 34,35 |
| 14.349 | DATA "ど", ${ }^{\text {d }}$ |
| 14350 | DATA "Z |
| 14360 | DATA , 4, 2, 30, 37 |
| 14.370 | DATA "K" |
| 14.389 | DRTA "E" |
| 14:399 | DATA , $2,4,40,43$ |
| 14408 | DATA $4,2,42,41$ |
| 14410 | CATA "M" |

```
14420 ORTR "\&", , ,
14430 DRTA "H",,.,
14440 DRTR "W"....
14450 DRTA ,3,1,45,51
14460 [ARTA , \(1,5,47,50\)
14470 DATR , 2, 4,4e, 49
14490 DRTR "Ү",,,
14490 DRTR "乡", , ,
14500 DATR "リ", 1,
14510 DRTR , \(1,5,52,53\)
14520 DATR "T", ,
14530 DRTA "F",...
```

If you are more confident（or are trying to save space）then all the DATA can be condensed on to eight rather unreadable lines which are OK for those who are good at counting commas．but very difficult to edit．

```
14019 [HTR , 1, 1, 2, 19, 1,5,3,19,
\(3,2,4,9,5,1,5,9,3,1,5,7, " 1,1\),
```




```
, \(5,7,19,15,, 2,4,17,16, " \mathrm{~A}^{\prime \prime}, \ldots, "\)
```



```
\(1,28,5,3,27,22\)
14220 UATR ,5, \(, 22,26,5,5,24,25\)
, "p", , , , "E", , , , "P", , , ,"L.", ,.
"["",.,., 5, 7, 45, 35, 2, 6, 31,44,
\(5,3,32,39,1,5,33,35,3,1,34,35\),
"必", ,, "2", ,, , 4, 2, 38, 37, "K"",
\(, \ldots " E ",,,, 2,4,45,43,4,2,42,41\)
, "阶"
```



```
\(\therefore, 3,1,46,51,1,5,47,56,2,4,40\)
, 49, "Y", ,,, "壮., ,., "U', , , , , 1,5
, 5'2,53, "T", , , "F", , .
```

To check the design produced against the patterns available（sec Flowchart 8．1），the array pointer $A P$ is first set to 1 so that the search is started from the beginning． X and Y coordinates are read from the $\mathrm{Cl}(\mathrm{AP})$ and $\mathrm{C} 2(\mathrm{AP})$ elements pointed to，and the last position LP pointer set equal to the current array pointer AP．

The point colour PC at these coordinates is determined by PEEK（TS + $\left.\mathrm{X}+\left(\mathrm{Y}^{*} 32\right)\right)$ ．If this is 175 then the point has been set and the＇yes＇pointer $\mathrm{Y}(\mathrm{AP})$ must be followed．If any other value is found then the＇no＇pointer


Flowchart 8.1 Character Recognition
$N(A P)$ is followed. In either case a check is made to see whether the element pointed to contains a zero (indicating the ultimate end of a branch), which shows that a character has been found. If so, the appropriate letter LE $\mathbb{( L P}$ ) is printed, and the display is held until a key is pressed, when a new cycle is initiated. As long as a higher value than zero is found then this must be another branch point and so the program loops back to 2010 and picks up the new values of $\mathrm{Cl}(\mathrm{AP})$ and $\mathrm{C} 2(\mathrm{AP})$.

To allow you to see which points have been checked, these are set to different colours as they are found. 'Yes' and 'no' branches can be distinguished, as tested points which were not set ( $P C=175$ ) will now be cyan $(239+16)$, whilst points which were not set will be magenta $(239+0)$. Any points which were set but not tested will remain blue.

## 2006 FP $=1$

```
2510 :
2@こG PC=PEEK(TS+X+(")
2G35 IF PC=175 THEN AP=Y(RP): TOT
0.850
2040 AP=N(RP)
20%01F RP=0 THEH 207G
2DED POKE TS+Y+CY*32)(239+C15*!
FC=175))\:COTO 251@
z9>0 PFINT LE$(LP);
2090 R&=INKEY&:IF R&="" THEN 298
\square
2BGD GOSIJB 130DG:GOTO 20
```

If you want to see which part of the tree was actually followed, then add these modifications which will print out the sequence followed a s a column of numbers to the right of the grid. The blankingstring BL\$ defined is used for partial screen clearance.

```
2005 PRINT { 16,"RP"::LI=1
2055 PFINT © ((L.I*32)+1E),RP L.I=
1.1+1
2995 FDF N=1 TO 10:PRINT R (CN*?
23+16%,EL.$:NEKT H/
13005 BL.%=" "
```

The disadvantage of this more rapid method, of only checking critical points, is that it will make a mistaken match if it encounters a shape that is not on the tree. whereas if all points are checked then no match will be found in such a case.

Early Optical Character Readers would only accept a single particular typeface, but the latest machines not only accept different styles of type, but actually learn the recognition rules for themselves by means of a built-in expert system. You teach these by showing them a few pages of text and then entering these same characters via the keyboard. However we feel that it will still be a long time before anyone can produce a machine that can read OUR handwriting!

## CHAPTER 9

## An Intelligent Teacher

Another place where Artificial Intelligence can be particularly useful is in teaching programs. It is all very well having a program which tests a student's knowledge at random, but this is not how a real human teacher works. As well as asking the questions, he keeps an eye ont he progress of the students, increases the difficulty of the questions as experience increases, and tests them more rigorously on the types of problems with which they are having difficulties. For example, if a child takes a test involving addition, subtraction, multiplication, and division, but only gets the division-type questions wrong, then it follows that the child should be given more division questions in the future to provide more practice.

Let's have a look at how we can introduce these "human' qualities into a teaching program.

## Questions and answers

We need to create random numbers to be used in the first question, which we will make addition. Using INT(RND(0)*10) will give numbers between 0 and 9 .



The computer adds these together and then goes on to an input and checking subroutine at 1000 .

## $40 \Gamma=A+B:$ ROSUE 100 O

First, the routine must print the question and input your answer IP.
1GOE FRIFT Fi; "九": E:"=":
101日 IFAPIJT IF

Your answer must then be checked. If the program answer $C$ is the same as your answer, then CORRECT is printed and the routine returns toline 40 . Otherwise WRONG is printed followed by the correct answer.

```
1029 IF C=IF THEH PFINT "OOFFEET
":PETIJRN
10%ด FFFINT "HPIONG, COFFRECT HIUSHE
F WhS ":
1F4E FETIJNH
```

The other three subjects（subtraction，multiplication，and division）can be easily dealt with in the same way if we replace the＇＋＇sign in line 1000 by a sign string SG\＄，which we can set to the appropriate character at the time． As $\operatorname{INT}\left(\mathrm{RND}(0)^{*} I 0\right)$ is common to all the calculations，we might as well define this as a function $R$ ．


$30 \mathrm{~B}=\mathrm{FH}$（ AD ）

5の $\mathrm{F}=\mathrm{FHR}(\mathrm{SIJ})$
6月 E＝FNP（SIJ）

80 A＝FNP（MI）
$90 \mathrm{E}=\mathrm{FNR}$ ！MII？

110 $\mathrm{A}=\mathrm{FF}$ 偪（ CI）
$120 \mathrm{E}=\mathrm{FH}$（R（DI）

1ロんD PRINT H：SGま：P：＂ご＂

Finally we jump back to line 20 to ask more questions．
140 LOTO 20

## Dividing by zero！

As it stands，the program cancrash if $B$ happens to bezero when a division is selected．This can be simply fixed by always adding one on to $B$ in this case：

1 その $\mathrm{B}=\mathrm{FHP}(\mathrm{CH} \mathrm{C})+1$

## Deleting decimals

We are using integer variables to keep us to round numbers，but of course a
division may still produce a fractional answer, eg:
$3 / 2=1.5$

To avoid producing decimals, A needs to bea multiple of B . To achieve this we calculate $B$ first and make $A$ equal to $B$ multiplied by a random number between 0 and 10 .

```
1106E=FPTP(01)+1
12W FA>INTC FHRCDI )
```


## Keeping a score

Now that we have the test itself working, we need to consider how to keep a score. The simplest thing is to increment a tries variable TR each time the subroutine at 1000 is used, and to increment a scorevariable SC each time a correct answer is obtained.

```
1010 IHFIJT IP:TR=TF+1
1M%G IF C=IF THEH:PRIHT "COPRECT
":SL=SC+1:GOTO 104O
10.4G PRINT "'חUוP SCGRE 1S ":SC:"
\prime";TF:FETINPN
```

If you prefer the score as a percentage then amend line 1040 as follows:

```
1046 PRINT "YCII HAVE HRC ":INT<<
SC/TF'洣10B ); "% CORFECT": FETIJFH
```


## How many questions?

As it stands the program will ask one question of each type in sequence, ad infinitum. We can limit the number by defining the number of questions NQ as a variable.

```
10 NO=32
```

Each time a question is asked, $N Q$ is decreased by 1 . and when $N Q=0$ the test ends (after eight questions of each type have been answered).

150 IF NQ>日 THEN 20
160 ENC
1010 INPUT IP $\cdot T R=T R+1: N O=N O-1$

## Shifting the emphasis

If we are going to bias the questions if favour of areas of difficulty，then we need to keep a record of performance in each individual area．We therefore need separate variables for each type of question（AD for addition，SU for subtraction，MU for multiplication，and DI for division）．These variables are defined in terms of one eighth of the total number of questions to be asked NQ．


```
FiO
```

Now if the correct answer $C$ is the same as your answer IP then an increment variable IN is set to－1，CORRECT is printed，and the routine returns．Otherwise IN is set to I，and WRONG is printed followed by the correct answer．

```
1020 IF E=IF THEH IF&=-1:PRINT "C
ORFECT": RETIIPN
103日 IH=1:FFINT "WFOHG, EDPRECT
AHSWER WASS "IC
1040 F:ETUF州
```

IN is added to the appropriate individual number of questions variable $\mathrm{AD}, \mathrm{SU}, \mathrm{MU}$ or DI on returning，producing an increase in this value if the answer was wrong，or a decrease if the answer was right．

```
40 SG=="+":C=A+E:GCSOU 1000: AO=A
O+IN
```



```
1.1+IH
```



```
m|lO+IN
```



```
[I+ IH
```

Now we add to check to see whether all the questions of a particular type have not been correctly answered（eg $\mathrm{AD}>0$ ，see $\mathbf{F}$ lowchart 9．1）．If all questions of one type have been correctly answered，then no more of this type will be asked as the line is jumped over．If the appropriate number of each type has been answered correctly（ $\mathrm{AD}=9, \mathrm{SU}=0, \mathrm{MU}=0, \mathrm{DI}=0$ ） then the program ends．

```
40 IF AD>ロ THEN SG紬+";に=R4&: &O
```

SUE $15900: R D=R D+I N$
 SLJE 1 日月D：SLJ $=S \mathrm{SIJ}+$ IH
 USIJE 1000：PIJ＝PMJ＋IN
130 IF DIン日 THEH SGぁ＝＂，＂：©＝ค P ：G DSIJB 10日G：DI $=[I+I H$
140 IF $R D=0$ RN［ $S U \because G$ RN［ $M I J=0$ Rrd ［ DI＝O THEN 160


Flowchart 9.1 Intelligent Teacher

Notice that you are no longer asked questions about areas in which you have correctly answered four questions without making any errors．If you make a mistake then AD．ctc．will be increased and so you will have to answer more than four correctly before AD reaches zero．

## Degrees of difficulty

How about making the questions easier or harder according to how well you are doing (ie the values of $\mathrm{AD}, \mathrm{SU}, \mathrm{MU}$, and DI )? So far the current values for $A$ and $B$ have always been between 0 and 9 as they were produced by $\operatorname{RND}(0) * 10$, but we now need to bias the numbers produced for the questions towards higher values, if you are correct, and lower values, if you are incorrect. At the same time, we must ensure that you do not produce negative values if your performance is abysmal.

The 'worst case' will be if you get all the questions right in three of the groups, and all the questions wrong in the last group. In this case only four questions will be asked on the first three groups, leaving $32-(3 * 4)=20$ questions to be asked on the last group. In addition we must remember that $X(e g A D)$ starts at a value of 4 , so that the maximum value of $X$ which could be obtained is $20+4=24$.

We therefore set up a weighting variable WT, which is calculated by subtracting three times the number of questions to be asked in each group (3*AD) from the total number of questions NQ and adding back on the number of questions in a group $A D$ at the start.

$$
\left.\mathrm{WT}=\mathrm{NQ}^{-}\left(3^{*} \mathrm{AD}\right)+\mathrm{AI}\right)
$$

This is more simply expressed as:

$$
\mathrm{WT}=\mathrm{NQ}^{-}\left(2^{*} \mathrm{AD}\right)
$$

$$
F[C \cdot b T=F I D-2 x+[0]
$$

We now replace the fixed value of ten by the difference betwen WT and X.

## 

To begin with, WT=24 and $X=4$ so numbers between 0 and 19 will be selected. If a correct answer is given, then $X$ will be reduced to 3 and numbers between 0 and 20 will be chosen. After four correct answers. $X$ will not change (for this type of question) as it will have reached zero and the line will be skipped. The last values will therefore be between 0 and 22 .

On the other hand if the first answer is incorrect then X will increase by 1 and the range of numbers produced reduced by $1(0-18)$. In the 'worst case' $X$ will be increased twenty times to 24 and (WT-X) will fall to zero for both A and B (so you should be able to solve that particular problem!).

## CHAPTER 10 <br> Putting It All Together

In the previous chapters we have dealt. from first principles, with various aspects of Artificial Intelligence. In this final chapter we have linked together many of these individual ideas in a single complete program.

The original 'intelligent' program was the famous 'ELIZA', which was a pseudo-psychiatrist program written to send up a particular style of psychiatric therapy. We have resisted the temptation to follow this lead and have opted instead to produce a replacement for the average computer salesman. This program combines some ideas on the processing of natura! language and on expert systems, to produce a result which should both understand your requests and make suggestions which take into account both your requirements and a number of hard commercial facts.

Enough words and values have already been included to make the program interesting, but you can easily customise it by adding your own ideas to the DATA. (We take no responsibility for the values included so far. which are for demonstration pur poses only, or for the views on particular machines expressed by the program!) The program itself is quite complex but it follows the methods described earlier in the book and the functions of the various line variables and arrays are given in Table 10.1.

## Making conversation

The format of the program is that you are asked for yourviews on each of a number of possible features in turn (the exact wording of the question being selected at random from a selection of phrases). Note that the key word or phrase is inserted into the sentence where necessary, and that the correct conjugation is applied.

Your input is examined in detail for keywords, and a rule array updated according to your requests. (If you want actually to watch the rule array being updated then delete line 5490 .) Many of the keywords are truncated so that one check can be made for a number of similar words. and a test is included to see if the matching string is at the start of a word.

The simplest answer is 'YES' or 'NO', which adds or subtracts I from the rule for that feature. If you mention the name of the feature (eg 'GRAPHICS') then a further $I$ is added to the rule. In addition, using a

Table 10.1 Main Variables in 'Salesman'

SIMPLE VARIABLES

| QP | no. of question sentences |
| :--- | :--- |
| Q | no. of questions |
| R | no. of rules |
| BB | bank balance |
| PH | phrase number |
| PHS | phrase words |
| M | match marker |
| OF | object flag <br> OM |
| object match |  |
| LD | like dislike |
| FS | rest of sentence pointer |
| NP | negative pointer |
| SI | AND match pointer |
| S2 | BUT match pointer |
| RU | rule update marker |
| OB | no. of objects |
| AJ | no. of adjectives |
| AV | no. of adverbs |
| L1 | no. of likes |
| DL | no. of dislikes |
| NJ | no. of negative adjectives |
| NV | no. of negative adverbs |
| HM | no. of cheap/expensive |
| C | no. of computers |
| FE | no. of features |
| CT | no. of cost ratings |
| CS | no. of cost suggestions |
| EX | no. of excuses |
| HI | no. of high price suggestions |
| LO | no. of low price suggestions |
| TC | total cost |
| TP | total profit |

## ARRAYS

| OB\$(OB) | objects |
| :--- | :--- |
| AJ\$(AJ) | adjectives |
| NJ\$(NJ) | negative addresses |
| AV\$(AV) | adverbs |
| NV\$(NV) | negative adverbs |
| L1\$(LI) | likes |
| DL\$(DL) | dislikes |
| Q\$(Q) | question objects |
| QP\$(QP) | question sentences |
| CR(Q) | cost rate |
| PR(Q) | profit rate |
| IC(Q) | total cost |
| IP(Q) | total profit |
| HM\$(HM) | cheap/expensive |
| R(R) | rules |
| CO\$(FE) | computer names |
| FE(CO,FE) | feature names |
| C(CT) | cost ratings |
| CS\$(CS) | cost suggestions |
| EX\$(EX) | excuses |
| HI\$(HI) | high messages |
| LO\$(LO) | low messages |

'positive' adjective or adverb adds to the rule, whilst a 'negative' adjective or adverb subtracts from the rule. Separating the words into different classes allows you to make more than one change to the rule at the same time

Thus:

YES

YES BASIC

YES BASIC NECESSARY

YES GOOD BASICNECESSARY

Whilst:

NO

NO MEMORY
adds one
adds two
adds three
adds four
subtracts one
subtracts two

Furthermore, verbs are grouped as 'likes' and 'dislikes', the last of which reverses the action of the rest of the words.

Thus:
I DETEST MACRODRIVES subtracts one
Both 'NO-' and 'N'T' are recognised, and most double negatives are interpreted correctly.

Thus:

## I DON'T LIKE SOUND

## I DON'T DISLIKE SOUND adds one

If anything appears at the start of a sentence and is followed by a comma, it is usually cut off and effectively ignored.

Thus:
NO, I DON'T WANT GOOD SOUND subtracts three

The exception is when 'AND' or 'BUT are included, when both parts of the sentence are acted on independently.

Thus if the question is:

## DO YOU WANT GRAPHICS?

and the answer is:

## NO, BUT I WANT GOOD SOUND

then one is subtracted from the graphics rule andt wois added to the sound rule.

If the program does not find any keywords in the input, it politely asks you to try again:

PARDON, EXCUSE ME BUT..
The program can only cope with one feature at a time, so if you try to ask for 'SOUND and GRAPHICS' at the same time, for example, you will get a request for a repeat of the question.

## HANG ON - ONE THING AT A TIME

However, it is possible to make comments about single features that you are not being asked about at the time, and these entries will still update the rules (as in the 'BUT example above).






## Decisions

In addition to the rule array, there are two other arrays which are linked to this. The first is the 'cost array', which gives an indication of the cost of this particular option, and the second is the 'profit array' which indicates to the salesman how much effort it is worth putting into selling this feature. The values for these last two arrays are produced by multiplying the content of the corresponding rule array element by factors entered originally as DATA in lines 10100 , etc, where the format is:
(phrase describing feature, cost, profit)

After each input, the salesman considers the consequences of your requests. First of all he looks to see if the sum total of the cost of all your requirements exceeds your bank balance. If so, he prints out one of a series of caustic comments on your credit-worthiness like:

## THIS SPECIFICATION SEEMS TO BE EXCEEDING YOUR CREDIT LIMIT

He also looks at howmuch profit he is likely to make on the sale sofar: if this drops too low, he will start to lose interest and come up with comments like:

I HAVE AN URGENT APPOINTMENT
or

## WE CLOSE IN FIVE MINUTES

At thesametime, hewill bemorehelpfulwith regard to which of the available computers will fit your requirements, drawing up a short-list by comparing the rating given originally to this feature in the description of each computer with the value you put on it. The format for the descriptions is:
(name, value of feature 1 , value of feature 2 , value of feature 3 , etc)

The highest rated machine will always be picked out first but, if possible, at least three machines (possibly with lower ratings) will be selected and the final choice is made from these. Either the highest or lowest cost computer (at random) will be selected for mention, for example:

IF YOU WANT A REAL ROLLS-ROYCE THEN JUST LOOK AT THE..
and
IF YOU ARE IN THE BUDGET MARKET THEN WHAT ABOUT THE．．．

If only one machine fits the bill，the program will come up with：
YOUR ONLY OPTION IS THE．．．

## Salesman

```
106 %OGIE G%N
```



```
F=IHSTFC1.FH!t, ",", IF SF %O THEN
```




```
LEH&FHक,-SF;
```



```
g" THERd F'H$:=L.EFT和 FH*,, SF-1 + "SS"
```






```
H市,LEN(PH& `-SF; ELSE FH隹=FH年+" "
+FTK
```








```
THEH IEG5
```




```
1.0以 IF SI+52=币 THEN 15以5
```






```
0):IF(O)=IF(O)+FR(O)
```




```
1765 IF SP% THEN FIJ:aFIJ+1:L[:=1:M
=1:ST=SP+1 GOTO 1601m
1OW5 SP=INSTF(ST,IN&, "NG")
195@ IF SP>的 THEN LD=-1:M=1:ST=S
F+1:NF=NFP+1: GOTO 1:555
2H5G SP=INGTF(ST,IN虫, "N"T")
2150. IF SF%以 THEN LLO=-1:M=1:ST=S
E+1:NP=FNP+1:GOTO 2SNG
20日月 IF NF%G THEN IF UNTCNP,O}=
F,2 THEN FIJ=FUI+1:L[:=1 ELSE FII=FIU
-1:L.[=-1
3Gn FOF N=币 TO L.I
2400 GF=IHETF(1.IN沾,II$(N):IF S
```



```
    THEN L[j=LD:\1:||=1
zrinf NE'\T N
```




```
F%@ THEN IF MIOक( IN惟,SP-1,1)=" "
    THEN L[~-LD*-1:M=1
##HyNTHTN
29#け FIJP N=ら TO CE
```



```
P%@ THEN IF MI[秋 INW,SP-1,1)=:" "
    THEHY PU=FOI+LD OF=N: M=1,OM
#10G HE%T H
zaらg FOF N=\squareSTO R\
305 SP=IHETF(1,IH直,FN$(H)):IFS
P-5 THEN 3505
```



```
EN 3605
20%5 FlJ=FlJ+t.C; :M=1
365G HENT H
3 7 9 0 ~ F O F ~ N H = [ J ~ T i J ~ N V ~
```



```
F=O THEN 41月G
3905 IF MIO&(IN&,SF-1,1)<>" " TH
EN 416JT
```



```
41BG HEXT H
4205 FOF N=0 TOLH.J
```



```
F-i.) THENH 4505
44951F MID$(1N怗,SF-1,1`<"" TH
```

```
E.N 46后S
```



```
4605J N位自 H
4T0G FONP NF=5
```



```
F=5 THEH& 5100
```



```
|f S1610
```



```
S16FJ NE:!T |
51t日 FDF H=&% TO Hy
```



```
F-5 THEF& %1SE
```



```
    5 1 9 5
```




```
E|OF}\mathrm{ IF SH=% THEH FF,1HT "FHTHEF
    ENFEHSINE"
519% NF%T H
```



```
F[HON. F'ERSE E SILE ME EIJT", SIITI]
    200%
```




```
TO !S0%%
```







```
4.459 B0T01%:%GF
5500% IL
```



```
Z2,FCHO, NENT H
```



```
1\Omega, 1C&H? NHEO H
```





```
ELGG}TE=TC+IC&
810% TF=TF+IFCN:
E2Gリ5 NE"T \
```





```
FFIHT:FRINT CSGEFT
```



```
GTEA FOR %=g TO Q STEF-1:POq=""
GSHG FON H=[ TO CO
60010 IF FE(N,0;-RC日)% THEN POW=
```



```
THDET HEST N
71DGO IF FIJ&="" THEN HE欺T %:GOTO
%205
#119 IF LEH&POW%S THEN HEYT K
7?M5 CLS
ア30G PRINT 巴 口.""
7316 G0T0%950
"35%% PRINT PO出
```



```
7%OD FOJR N=1 TO L.EN& PO%)
TFGG PPINT CO$(UFL(MIO$(PO$,N,1)
    ) \
7THO NEMT N
P9505 PRINT
7コロด T:ラ=ด: E'5=16
```



```
B105 NG=\/RLMIC${PO&,IH+1,1)%
SOD IF C(NC%)=TS THEN TS=C(ND):
HI=1绾
3%Gg IF C(HC)<=ES THEN BG=LC(NC):
LO=NC
O405 ME<T CH
8410 IF HI=LOJ THEH FRINT"'%OUJR OHA
L.Y OPTION IS THE" FRIMT CO官HI):
%JTO Э20以
350G HI*=LD%(HI):LO&=CO$(LO)
S5gh SE=RNO<2)
8700 SL=R4[43)-1
9305 IF SE=2 THEN 9105
395D PRINT HI${SL. %,HI$
gg@F GOTO 920%
9100 FRINT L.O$(SL ),LOU
0209 6-0+1:IF O20 THEN 20, ELSE
    ENND
9:305 OP=C5: 
5:LI=S:[L=3:NJ=O:N|=2:HIN=3:[IM O
```





```
3,H|IWE HM Y:EE=1月ด
```




```
IMFTWE, ISN, SOFTWHFE,FHFTFIDGE, I
OGTILR, FSGENEL, FENTFOHIE, FOOZG,
ESFH!LD, HETLUNF, 15-EIT,MIL.T ITHGF,
SEF"\IDE
```



```
F,FIFGT, FAST, EFFLE, ESGEHT, LOT
GBOG OFTH EHLH, FIJEE[SH, FODF,GLIOW,
IHEFFIG, FEW, NOPG, LEHET, LESS
```



```
ECEEN. TFU
```



```
###, [PTTH WH怆, IKE, HFE[;,FEO|IIFE
1WHFFH [HTH HATE,OISL IFE, LIOHTHE, D
ETEST
```





```
AFHE MEMOP%, 
```



```
&ENTETSIME EOFTWHFE, G, G, #H CHF
TFI[NE FOFT, 1.G
IGOG [HTH :H ID'GTICK FOFT, 1,\vec{F}
```






```
5,%合哈 GEPWIS:E,1,9
1FSGF] [HTH WOlllO MDII L.TFE, WHHT F
```



```
U FECJIRE, & IMFOFTHKT
1GO1日 [FTH EHENF,INE;隹ENGI',E
1G?% [JHTH [EHF, EVPEHGIVE
```



```
HENT H
```



```
HE:T H
```



```
HE%T H
```

```
1070% FOJF N=5 TO F%:REHE F%'क(N):
HEYT H
```



```
HE% H
10G[HA FOF HOE TO L.I:FEFE L.IDCN
HE'`T H
```



```
HENT H
```



```
(Hり,PFCN) HEXT H
11200 FOF HFG TG OF PEFGN OP&CN?:
HE`T H
```



```
HENT H
11290 ULS:9=0
11400 FFIHT "IT 15 M'' FLEAGUFE T
```



```
ILFOETOPE"
1150. FFINT"WE HFE INHEOHETE['L'' T
HE IILTIMATE SOUFEE DF FLL DOMFUT
EF PFODOLGTG HHN J SHALL HANE GPE
HT FLEASUNE IHA HELF IHG 'HOLISELES
T ODUF r, MEN MELHIHE:"
11GD日 FFINT"SO THAT I EH& WOFK O
IIT THE EEST CD|PITER FOF 'IOIF FA
ETICLILAR HEEOS FERHAFS 'ROUN WO
|L[ EE KIHE EHOUNH TO FHEWEF A F
EW THESTIONS"
```



```
11900 00=9:FE=19 T=O:TM [01$(FE
3,FEOD,FE, DF(CD,FE), COCT)
```



```
0, -, , 7, 7, 0, 7, E, 0, , 0, 9,9
1,2019 [HATH FHALT SEFIDNE,6., , 6,8
```



```
1210. DHTH LLEHFSIM MT,9,9,9,7,7
,9,0,G,9,6,\vec{י},\vec{0},\vec{0},6,\vec{7},9,9,9,1
12己Q@ [HTA FHHROH ILLIISIOH, *, 
```



```
*
123%E [:ATA EHNANA IIE, 3,5,2,5,0,
4,6,日, 2,0,2,0,0,0,6,7,0,0,0,4
```




```
LHSS FROOUNT THEN 'FOU PMET
    TFY' THE FGIF STHTE DF THE RFT TE
```



```
DII WAPHT A BOLLSGFDGEE THEH JIST
    LOOKK HT THE
```



```
(.).)
```



```
HEYT H
14904 FOJP HF& TO HI:PEFC HIWCH?:
HEMT H
14SDA ELG:RETURH
```


## Commentary

Lines 200-440: Pick the words to be used in the next question, and select the correct conjugation.

Lines 500-800: Set up your INPUT and reset variables.
I ine 900: Checks for a comma.

Lines 1000-1200: Check for 'AND' and 'BUT'. If neither of these is present the program jumps to line 1500 .

Line 1300: Updates the current rule negatively if 'AND' or 'BUT' are present and the first word is ' NO '.

Line 1400: Updates the current rule positively if ‘AND' or 'BUT' are present and the first word is not ' NO '.

Line 1500: Deletes anything preceding a comma.

Lines 1600-2100: Check for 'YES', 'NO' and 'NT' and update the current rule accordingly.

Line 2200: Checks for a double negative.

Iines 2300-2500: Check for 'likes'.
I.ines 2600-2800: Check for 'dislikes'.
I.ines 2900-5100: Similarly check for objects, adjectives and adverbs.

Lines 5110-5190: Check matches for high and low cost key words.
Line 5200: Checks for no match and reports.
Line 5300: Checks for more than one object.
Line 5400: Updates the current rule, or another rule, accordingto whether or not the object matches the current question.

Line 5490: Jumps over the print-out of the rules.
Lines 5500-5800: Print out the rules.
Lines 5900-6200: Update the total cost and total profit values.
Line 6300: Prints an excuse if the profit seems too low.
Line 6400: Prints a warning if the spending is too high.
Line 6500: Zeros the total cost and profit values.
Lines 6700-7120: Search for computers which match your requirements.
Line 7310: Jumps over the print-out of matching machines.
Lines 7350 - 7800: Print out the matches.

Lines 7900-8400: Pick the highest and lowest priced machines which match the specification.

Line 8140: Checks if only one machine was selected.
Lines 8500-9100: Print out the name of either the highest or lowest priced machine.

Line 9200: Updates the feature to be checked and returns for another input.
Lines 9300-11300: Enter the information on features, keywords, costs and profits.

Lines 11400-11700: Provide an introduction.

Lines 11800-13800: Enter the information on the names and virtues of particular machines.

Lines 13900-14300: Provide warnings and excuses.

Lines 14400-14900: Contain the words $f$ or high and low cost messages.

## The rest is up to you

Artificial Intelligence is a fascinating subject, and we trust that we have given you enough information to get you started on your own experiments in this area. We have certainly enjoyed making our own explorations whilst putting this book together, but we have started to wonder how long it will be before someone designs an expert system program which writes books
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Artificial Intelligence on the Dragon computer shows you how to implement Al routines on your home micro and turn it into an intelligent machine which can hold a conversation with you, give you rational advice, learn from you (and teach you) and even write programs for you.

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Keith and Steven Brain are a father and son team and have already published the best selling Dragon 32 Games Master and Advanced Sound and Graphics for the Dragon computer. They are both regular contributors to Popular Computing Weekly.


[^0]:     1 ODGD DATR NORTH,NDRTH, $日,-1$, SIUT H, SOIJTH, $\sigma, 1$, WEST, WEST, - 1, $\quad$, ERST, ERST, 1, $\square$

